

TFNP:

- Collapses
- Separations
- Characterisations

Goos, Hollender, Jain, Magstre, Pires, Robere, Tao

**TFNP** = Total Function NP [Papa 80s]

Polgtime  $R(x, y)$

Input:

$x$

Output:

$y: R(x, y) = 1$  &  $|y| \leq |x|^{O(1)}$

Promise:

$R$  is TOTAL  $\forall x \exists y: R(x, y)$

# Sink-of-DAG

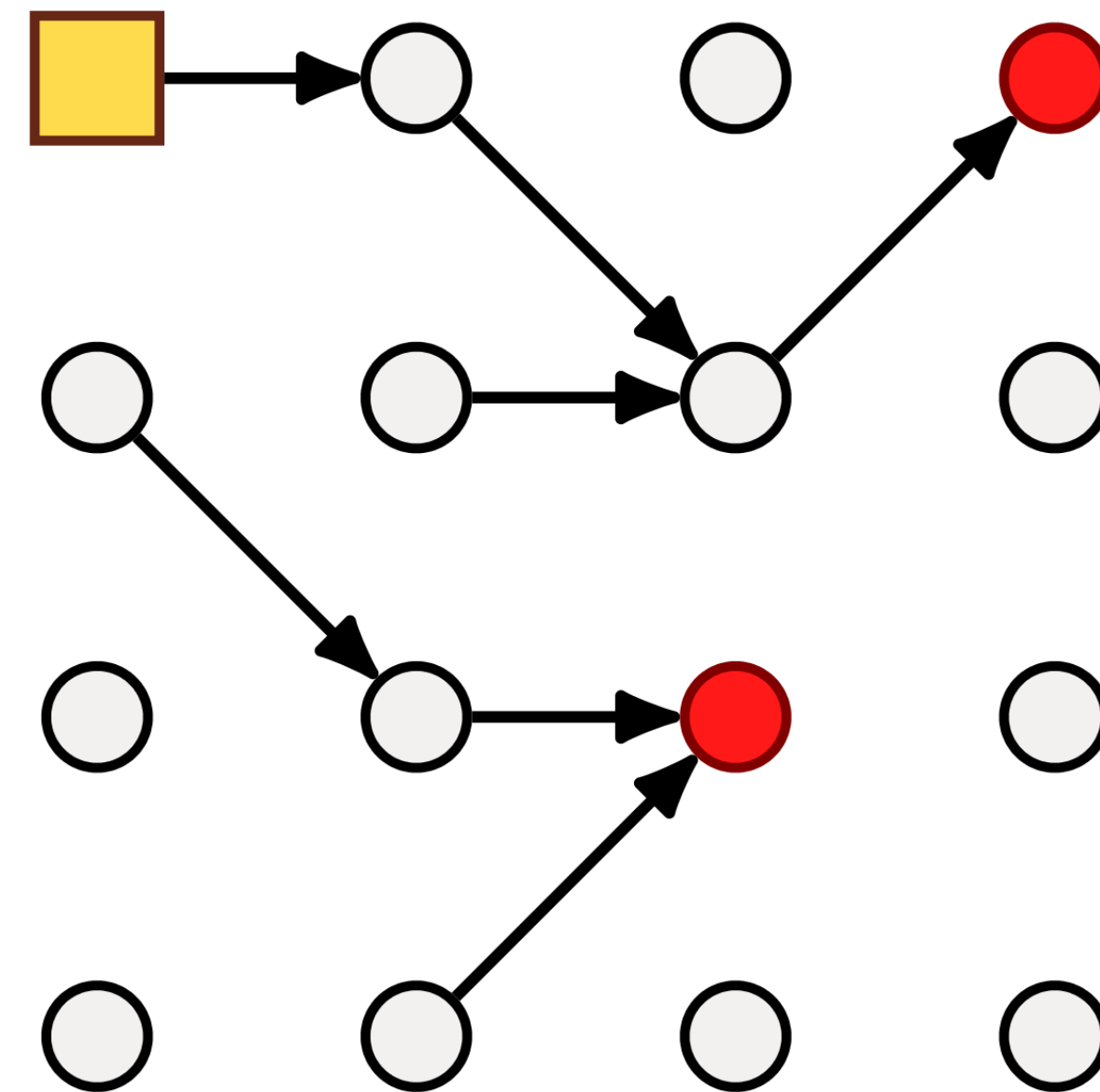
input

Circuit  $S: \{0,1\}^n \times \{0,1\}^n \rightarrow \{0,1\}^n \cup \{\perp\}$   
Defines DAG on grid  $[N] \times [N]$ ,  $N = 2^n$

output

sink node 

(or  $(1,1)$  is  $\perp$ )



Def: PLS = {  $P \in TFNP$  :  $P \leq SoD$  }

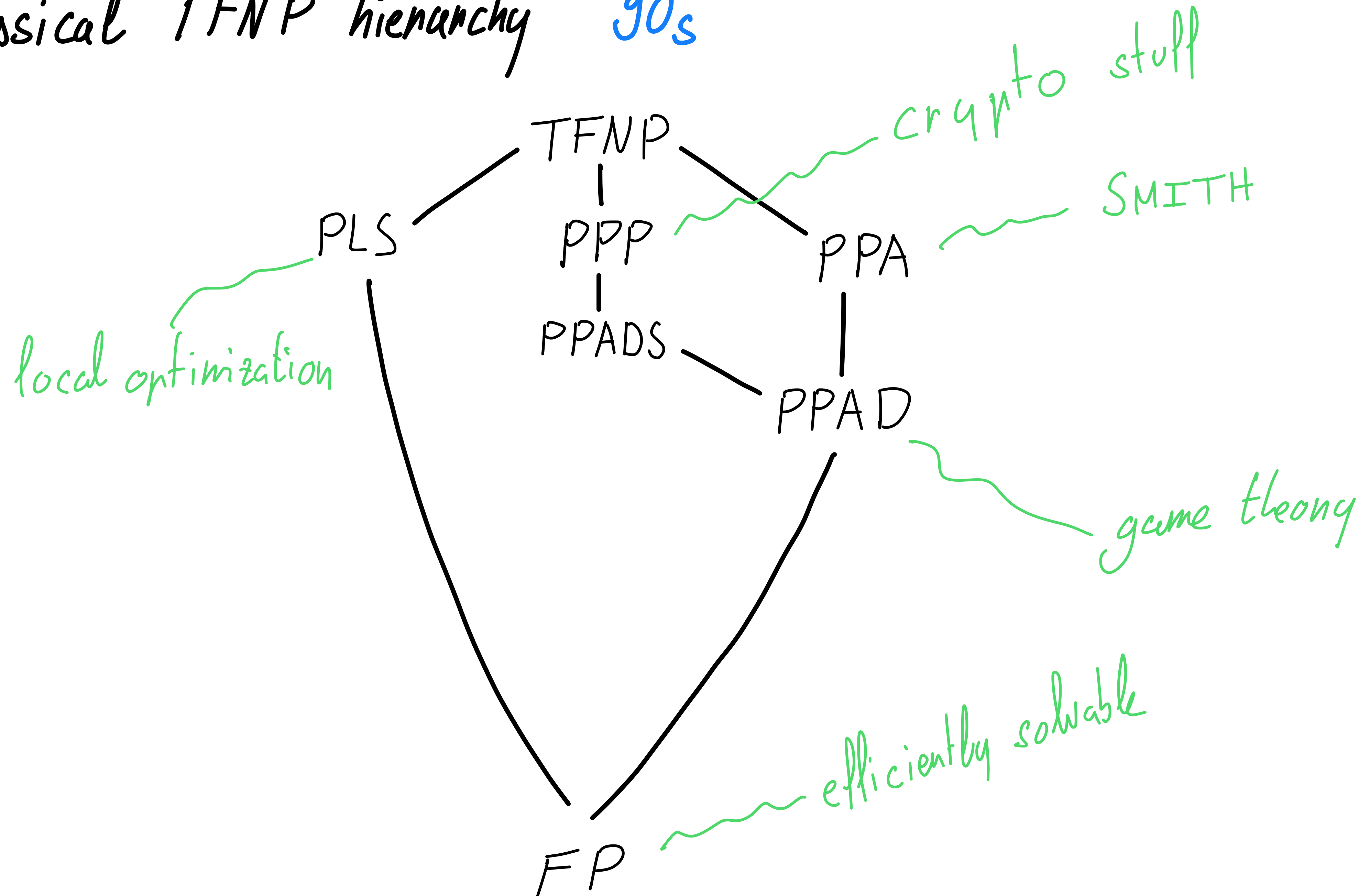
$A \leq B$  iff  $\exists$  polytime  $f, g$

input  $x$  to  $A \xrightarrow{f}$   $f(x)$  to  $B$

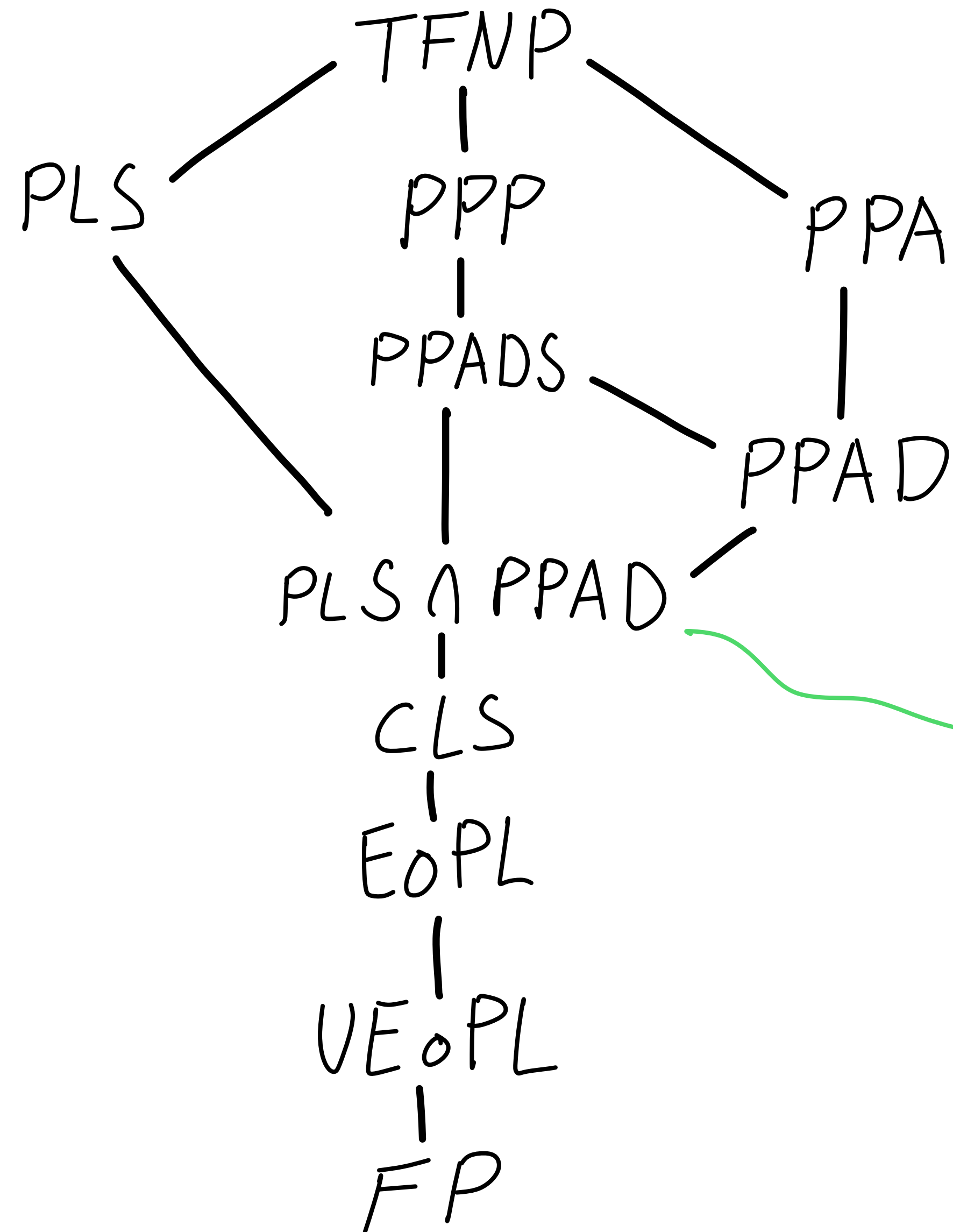
Alg for  $B$

$A(x, g(x, y)) = 1 \xleftarrow{g} y : B(f(x), y) = 1$

# The classical TFNP hierarchy 90s



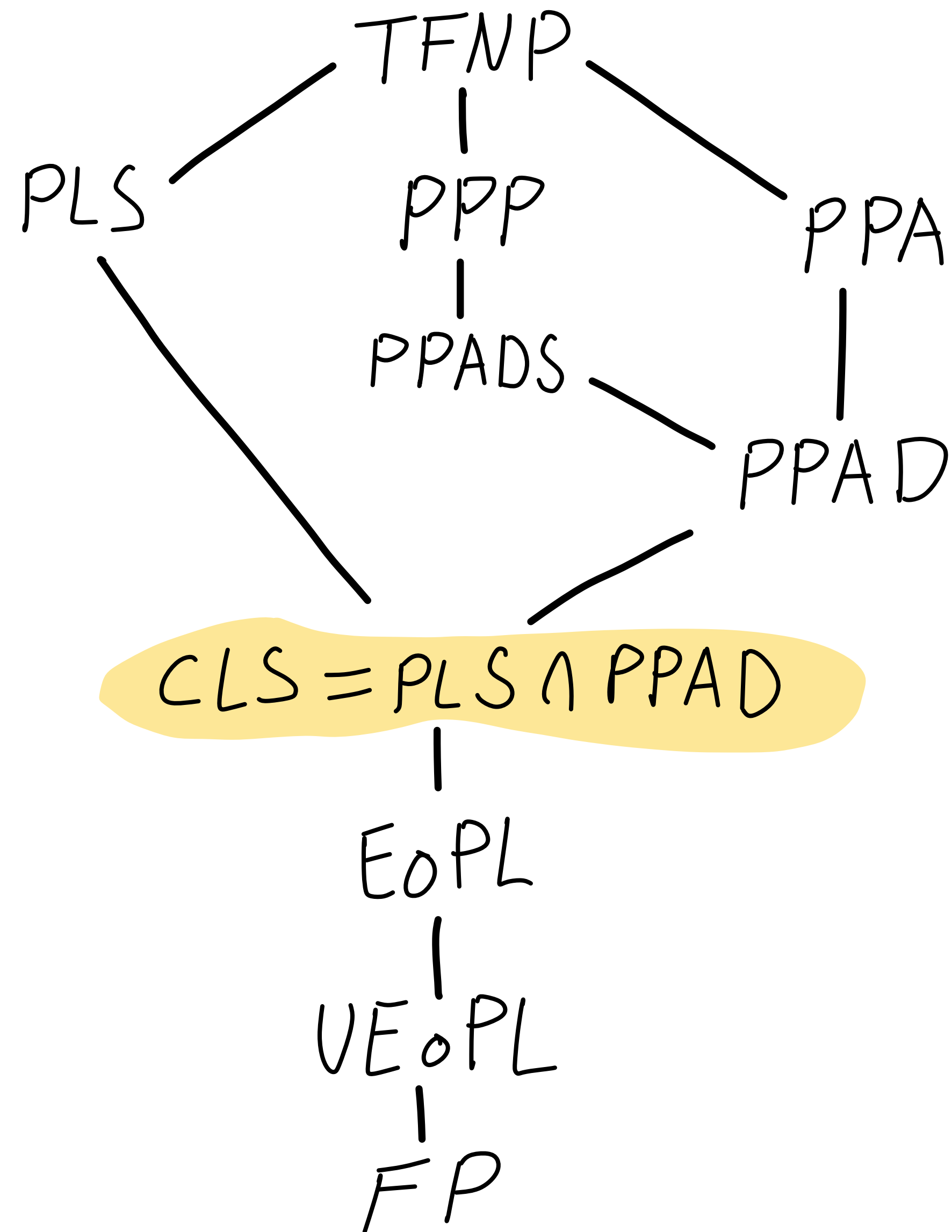
# New lower classes (10s)



un-natural class that contains many interesting problems

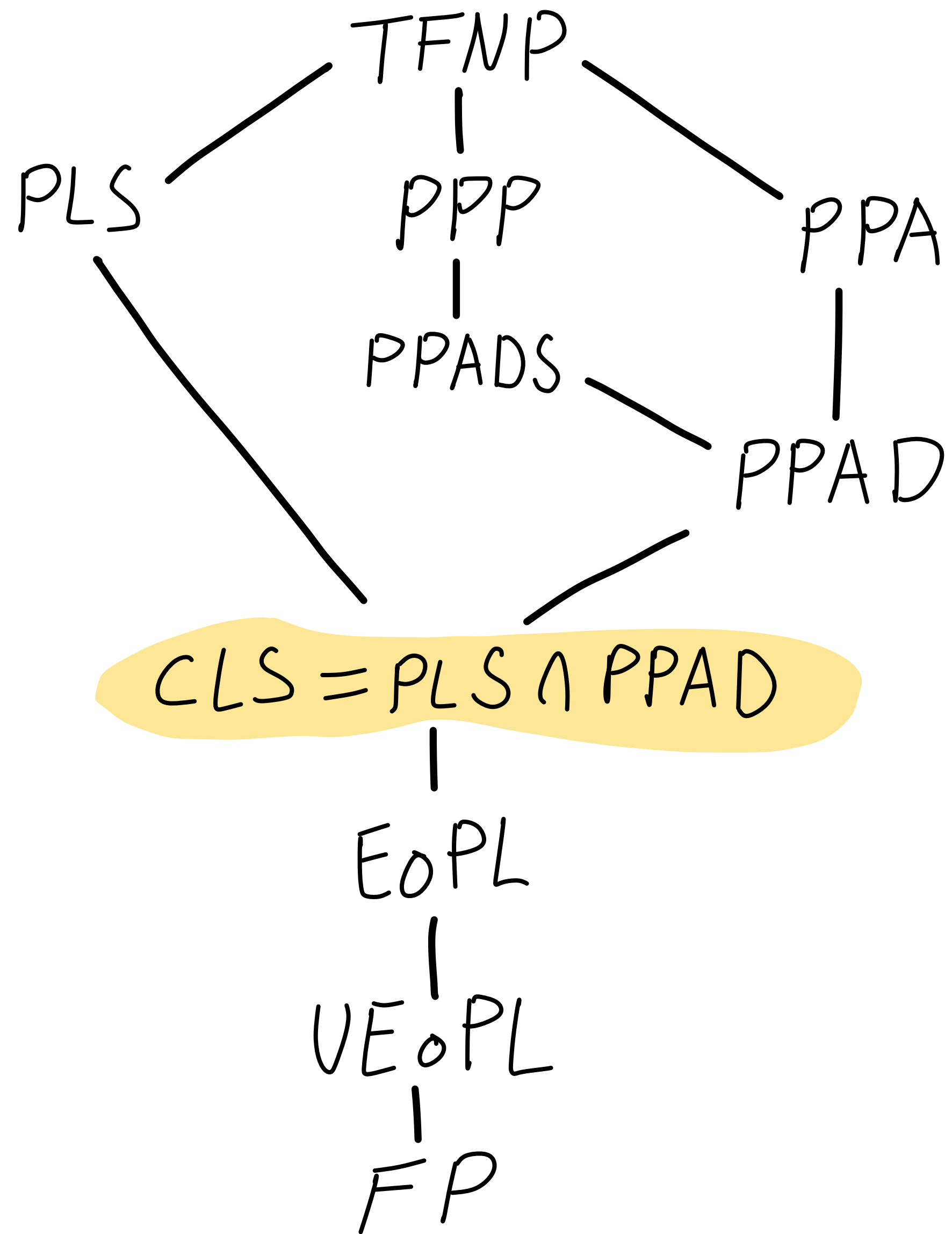
# A breakthrough collapse (2020)

[FGHS 20']  
Best Paper!



# A breakthrough collapse (2020)

[FGHS 20']



1. Are there any other collapses?

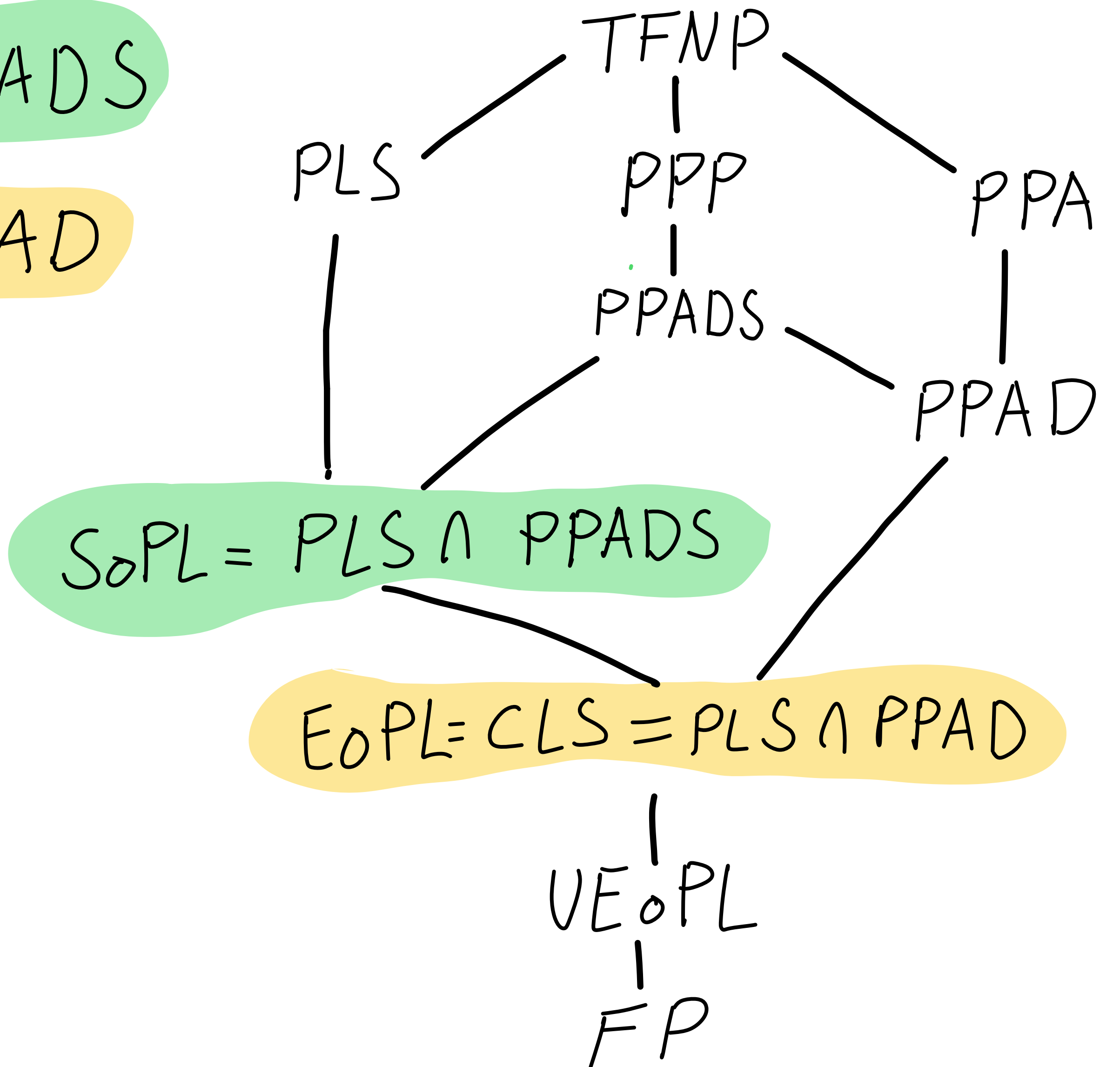
2. Can we argue against collapses?



# Result I: two further collapses

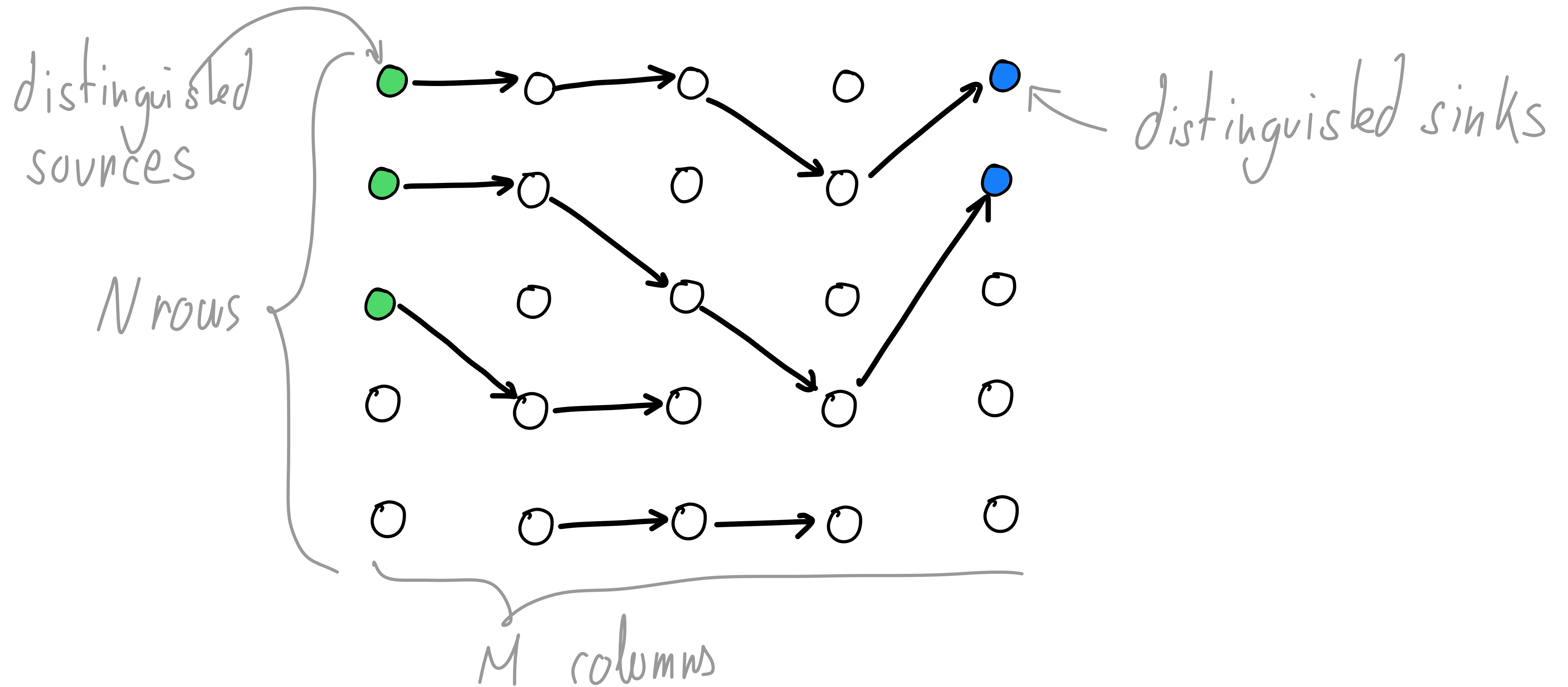
Theorem 1:  $S_0PL = PLS \cap PPADS$

Theorem 2:  $E_0PL = PLS \cap PPAD$



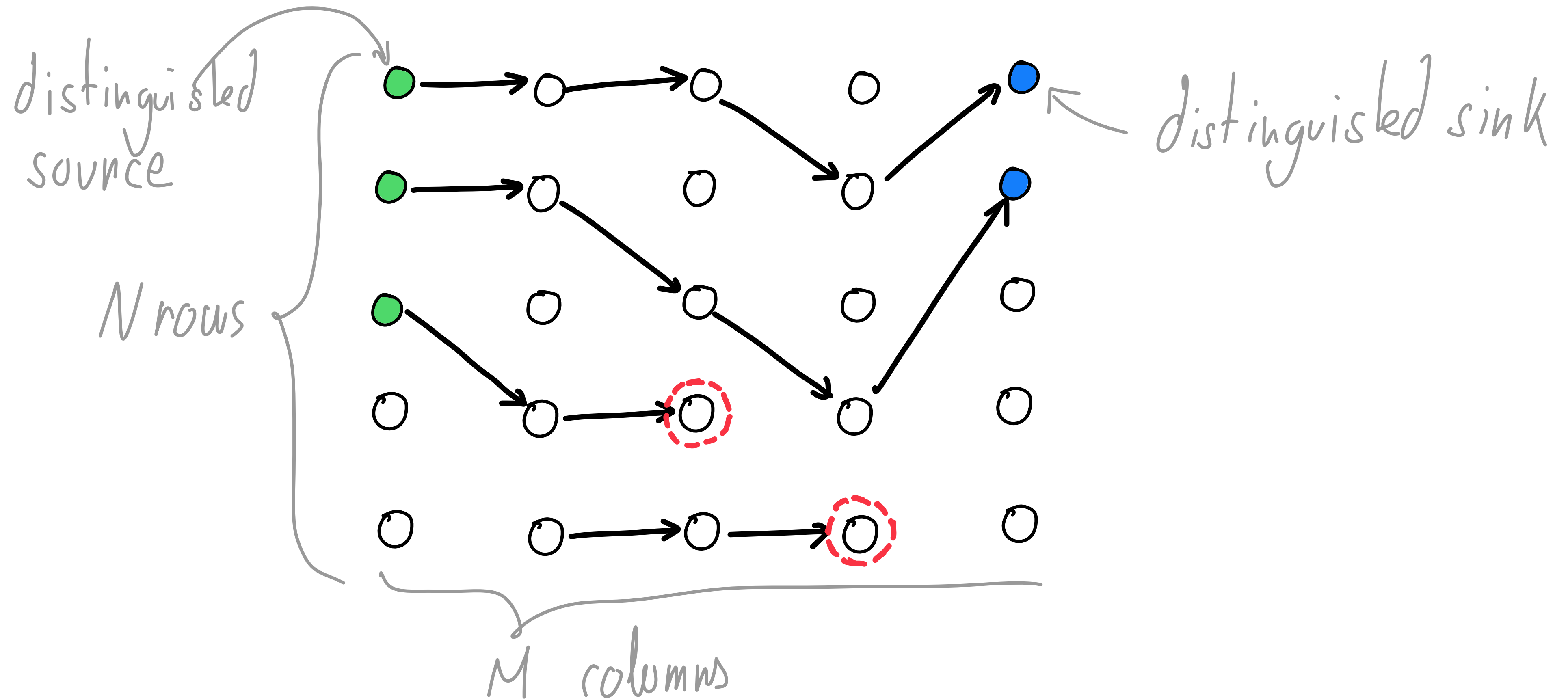
More on **EoPL = PLS  $\cap$  PPAD**

Idea 1: a grid view of TFNP



More on **EoPL = PLS  $\cap$  PPAD**

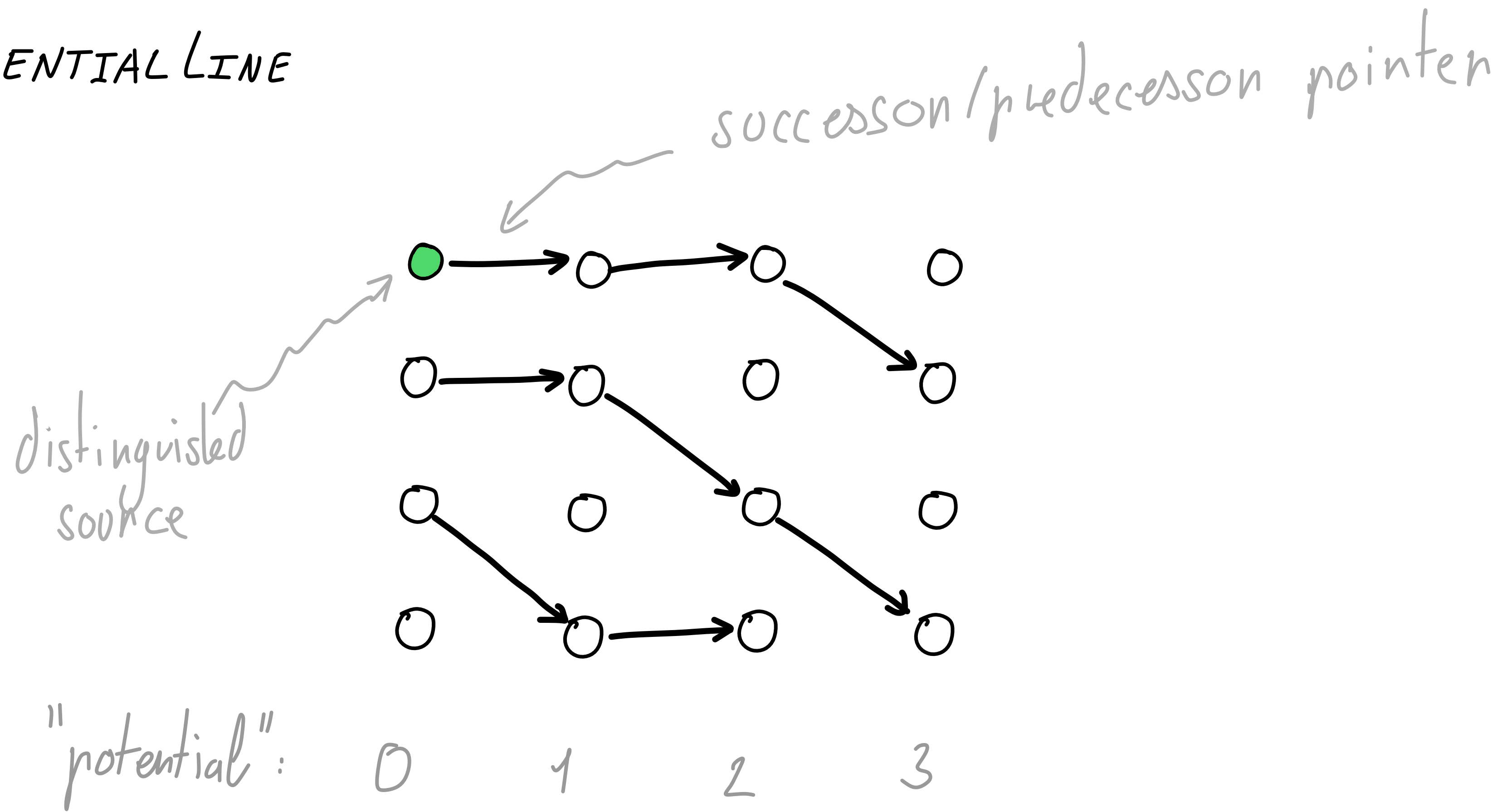
Idea 1: a grid view of TFNP



More on **EoPL = PLS  $\cap$  PPAD**

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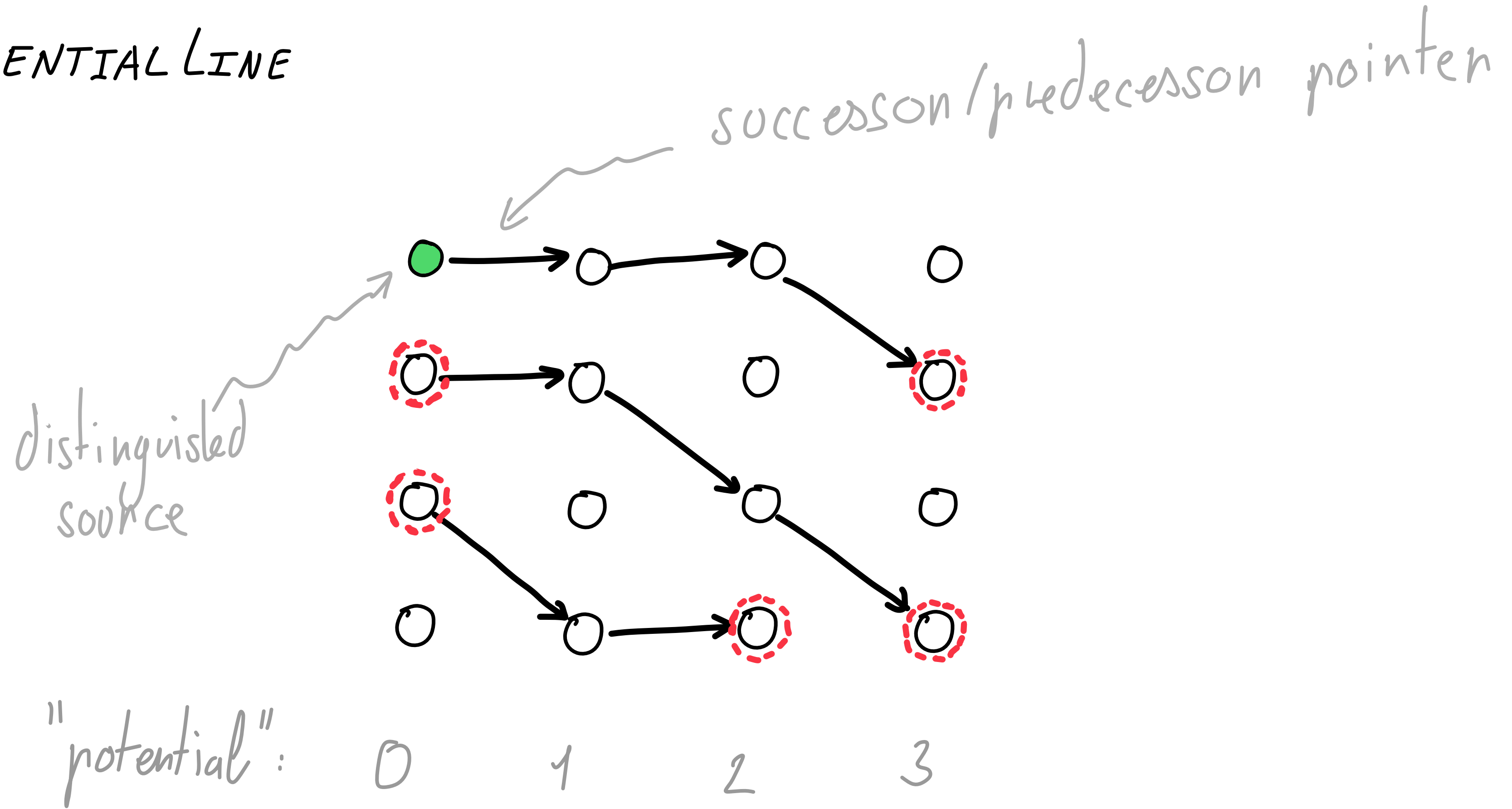
EoPL : END OF POTENTIAL LINE



More on **EoPL = PLS  $\cap$  PPAD**

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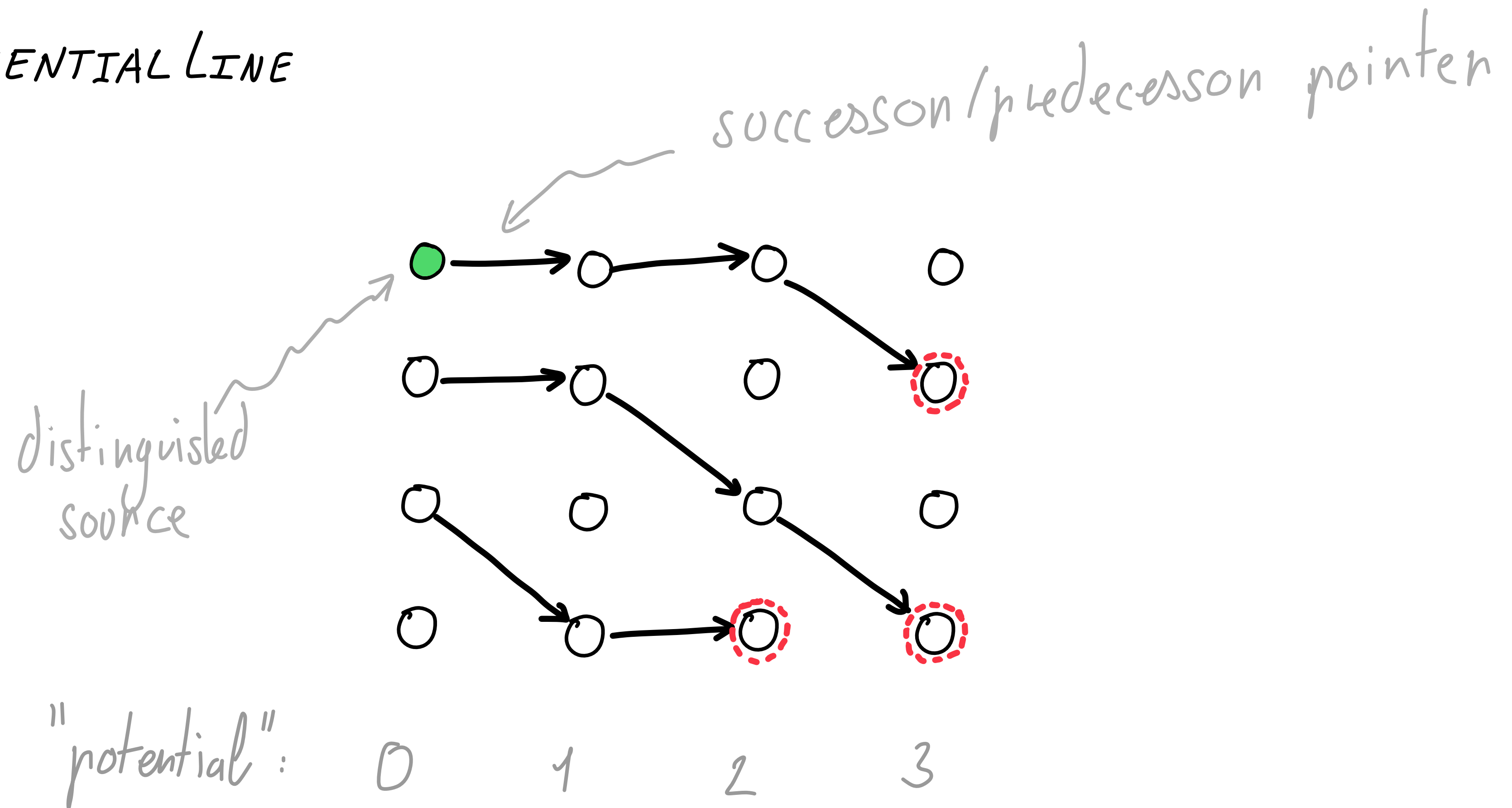
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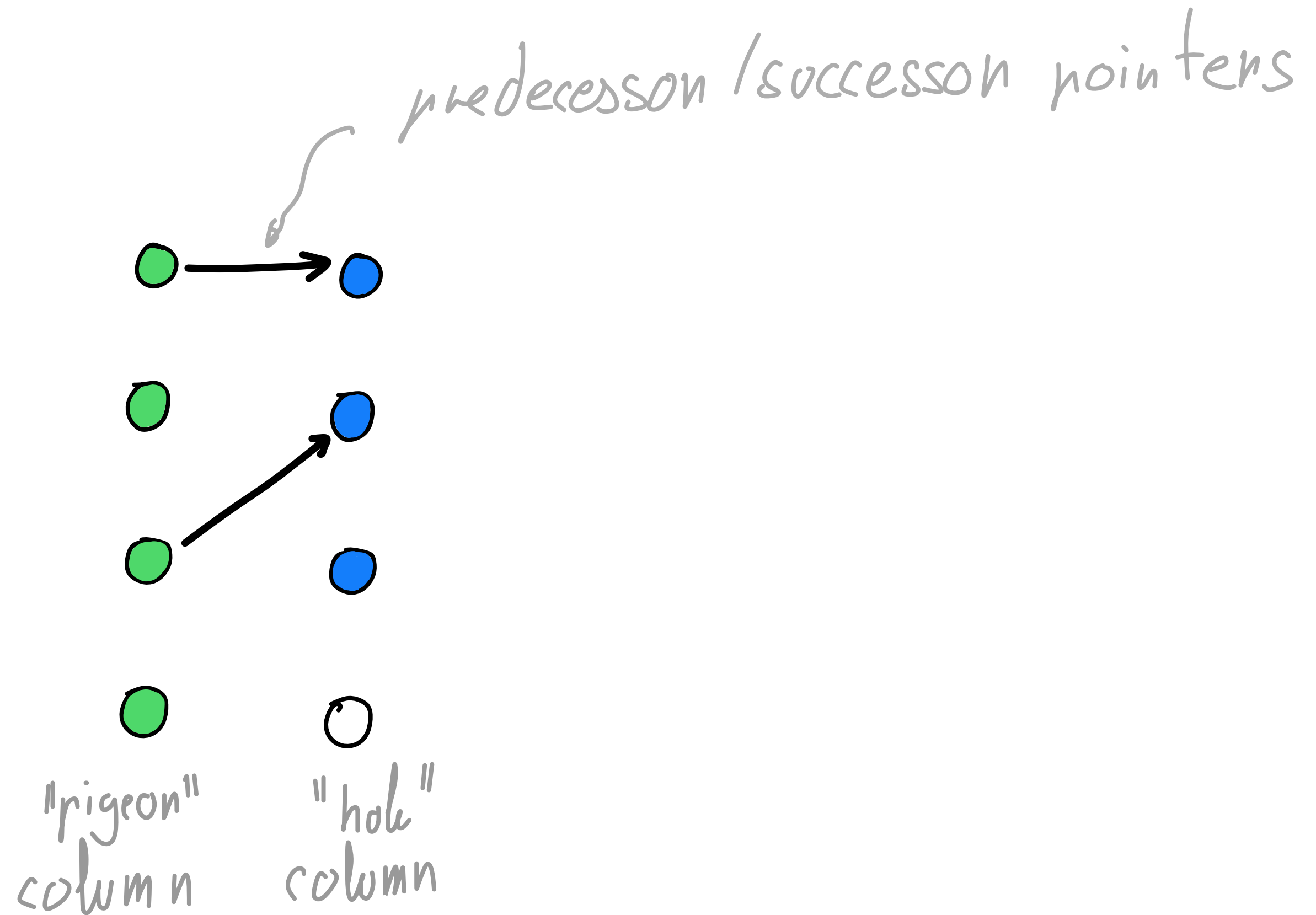
SoPL: SINK OF POTENTIAL LINE



More on  $E_0PL = PLS \cap PPAD$

Idea 4: a grid view of TFNP

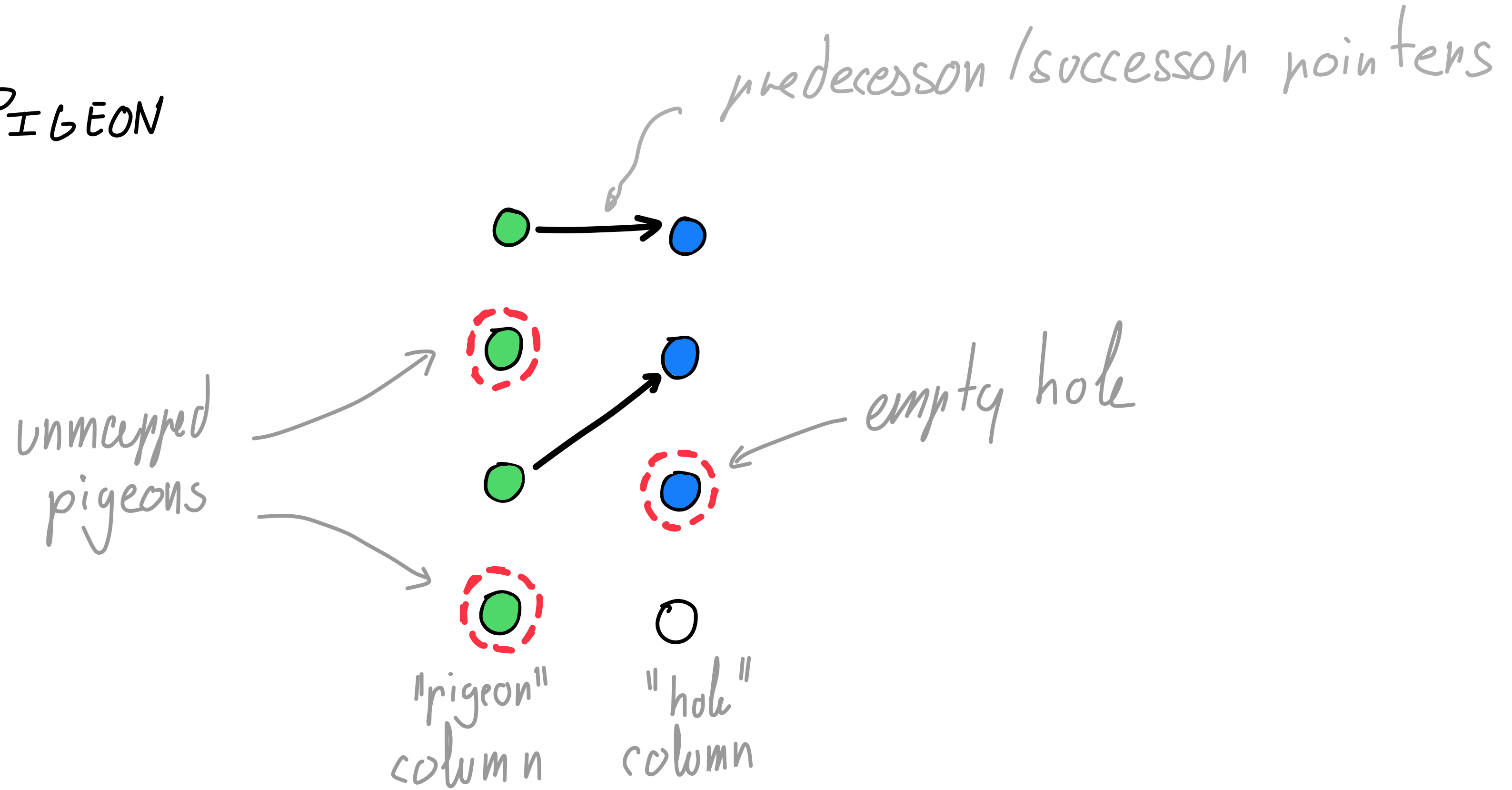
PPAD: BIJECTIVE PIGEON



More on  **$E_0PL = PLS \cap PPAD$**

Idea 4: a grid view of TFNP

PPAD: BIJECTIVE PIGEON





More on  $E_0PL = PLS \cap PPAD$

Idea 2: use  $S_0PL = PLS \cap PPADS$

- Only need  $PLS \cap PPAD \subseteq E_0PL$
- enough:  $S_0PL \cap PPAD \subseteq E_0PL$
- complete problem for  $S_0PL \cap PPAD$ :

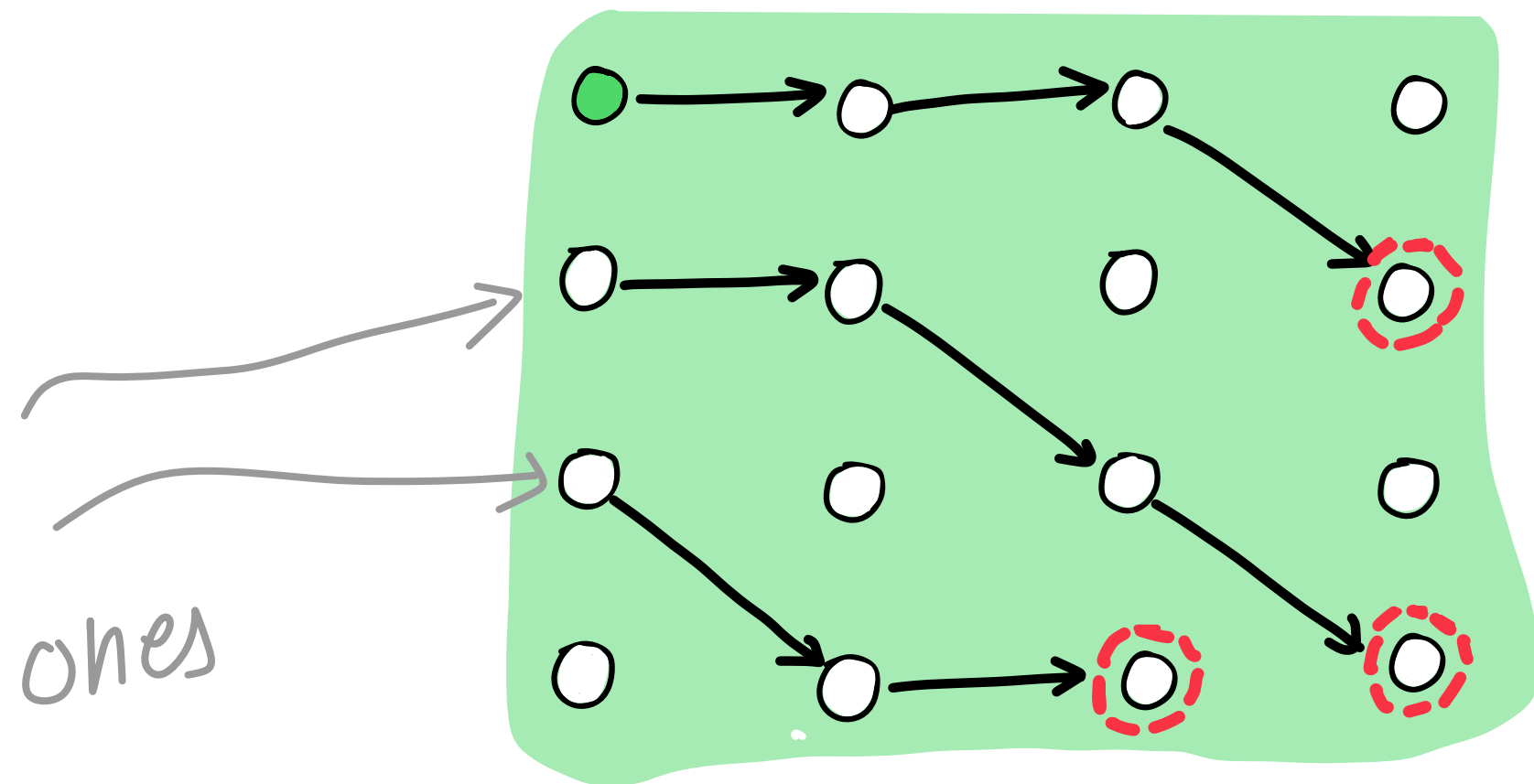
EITHER (SINK OF POTENTIAL LINE, BIJECTIVE PIGEON)

↖ receive two instances, solve any of them

# More on $EoPL = PLS \cap PPAD$

- Only need EITHER ( $SoPL$ ,  $B_{\Sigma}$  Pigeon)  $\leq$  END OF POTENTIAL LINE
- can we do  $SoPL \leq$  END OF POTENTIAL LINE?

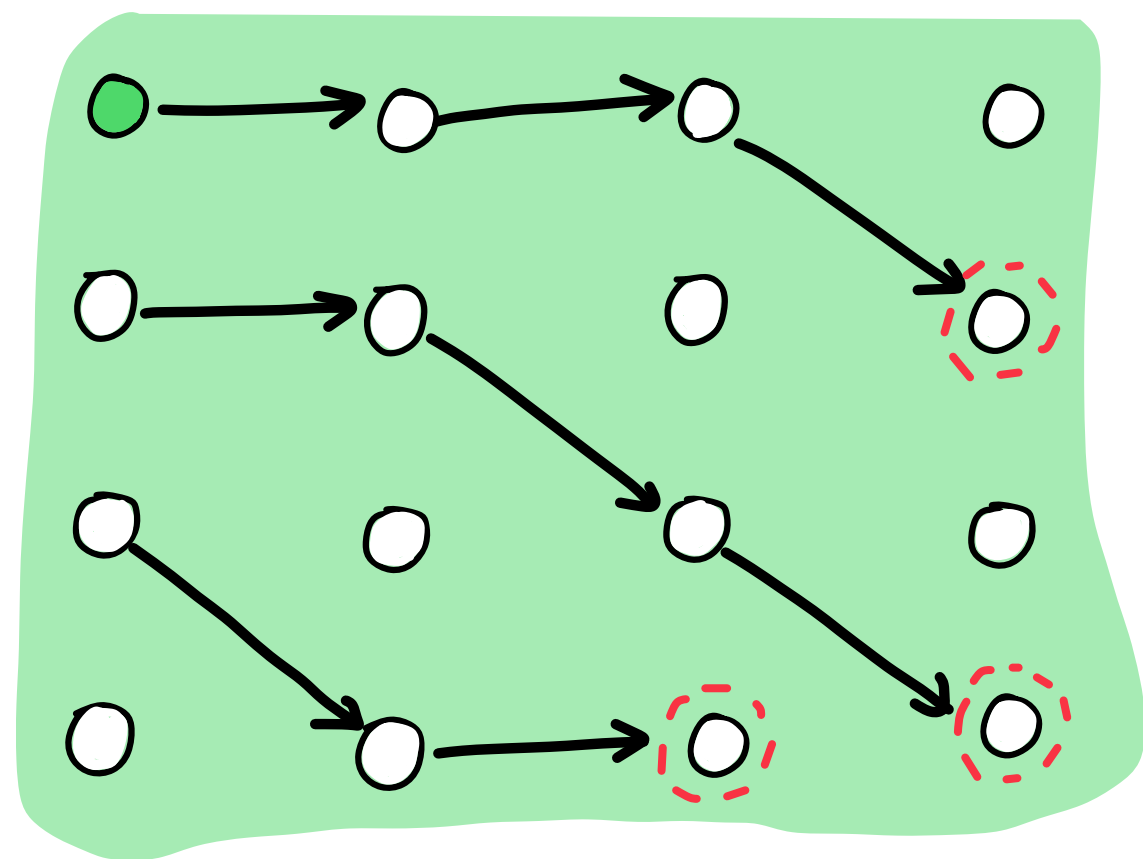
Sources are  
EoPL solution  
but not SoPL ones



$SoPL$

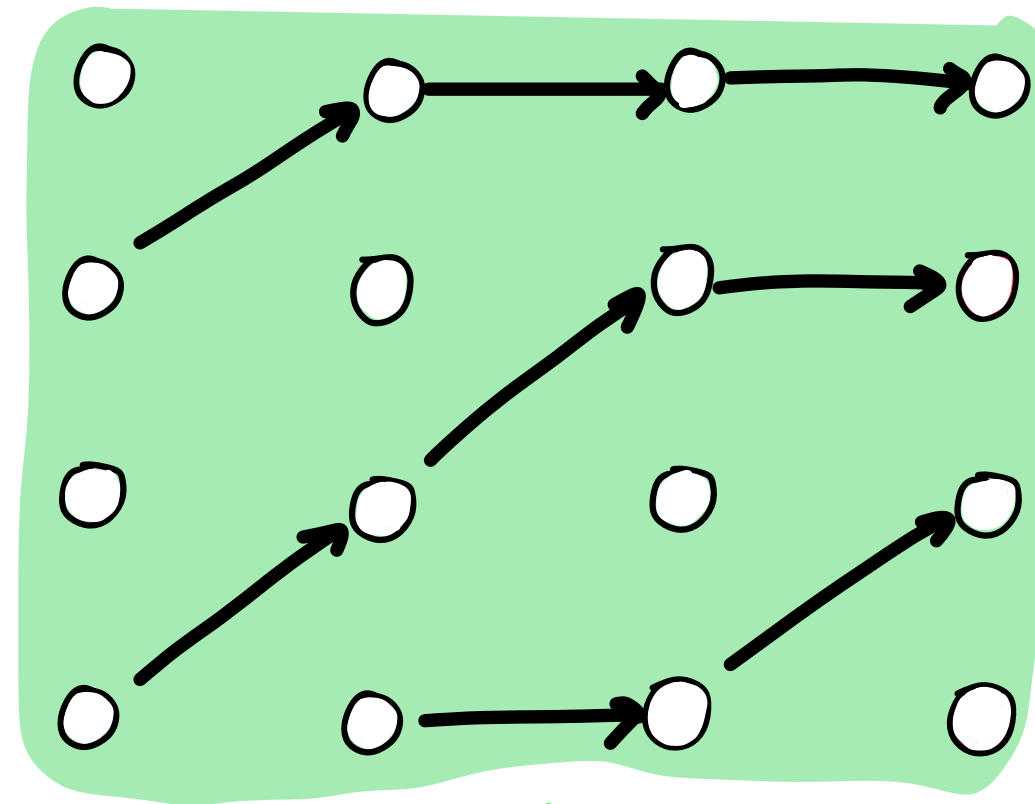
# More on $E_{OPL} = PLS \cap PPAD$

- Only need EITHER ( $SoPL$ ,  $B_{\Sigma}$  Pigeon)  $\leq$  END OF POTENTIAL LINE
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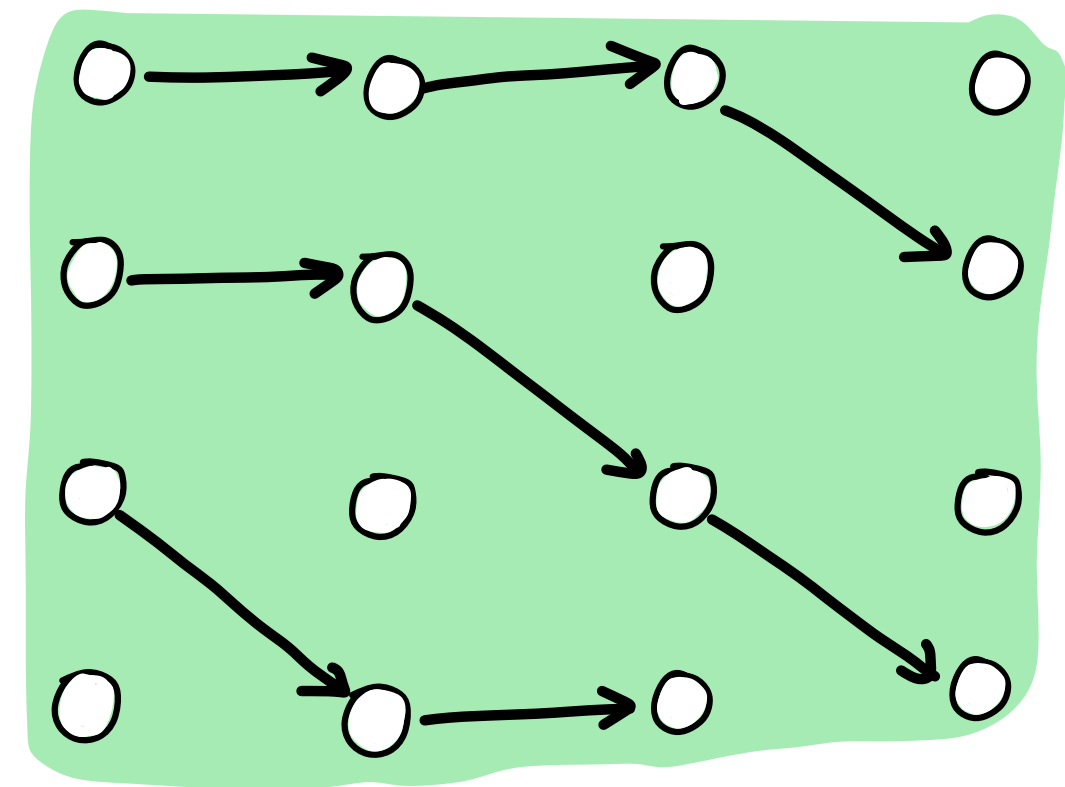


$SoPL$

or



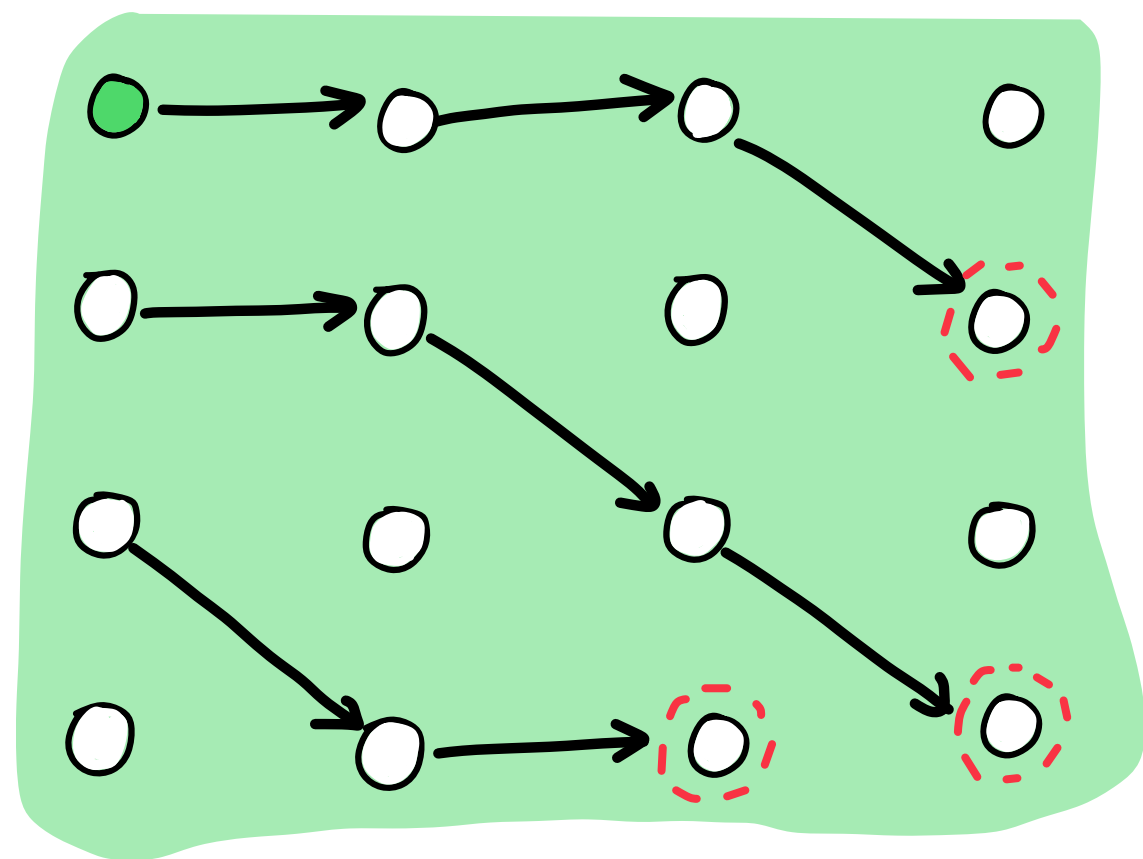
reversed  $SoPL$



$SoPL$

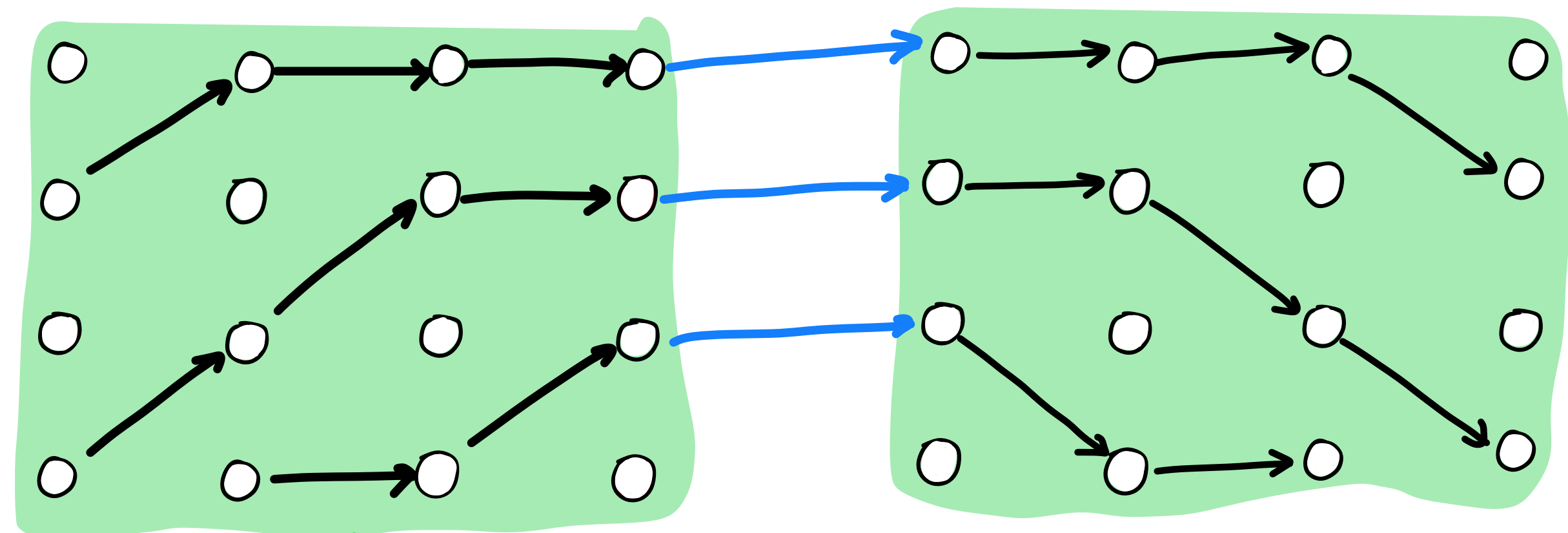
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- Only need EITHER ( $SoPL$ ,  $B_{\Sigma}$  Pigeon)  $\leq$  END OF POTENTIAL LINE
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$SoPL$

or ?



reversed  $SoPL$

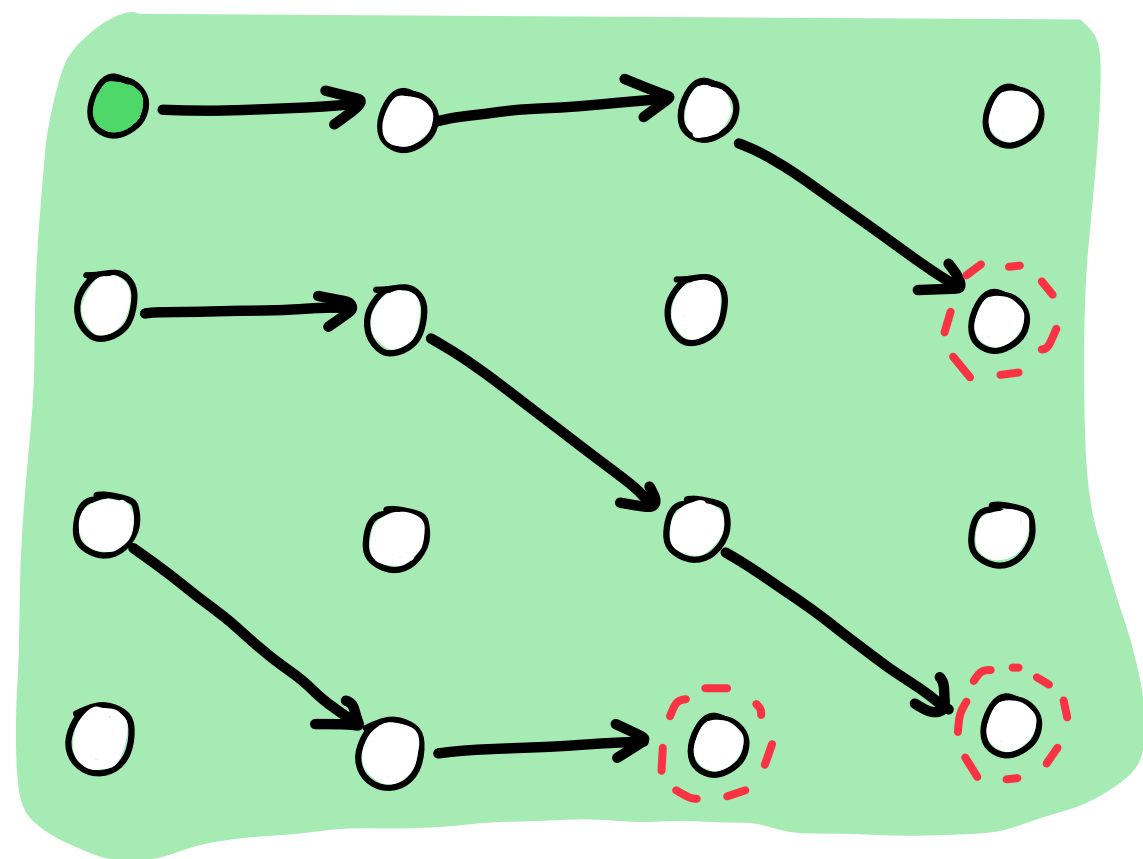
$SoPL$

END OF POTENTIAL LINE

# More on $E_{OPL} = PLS \cap PPAD$

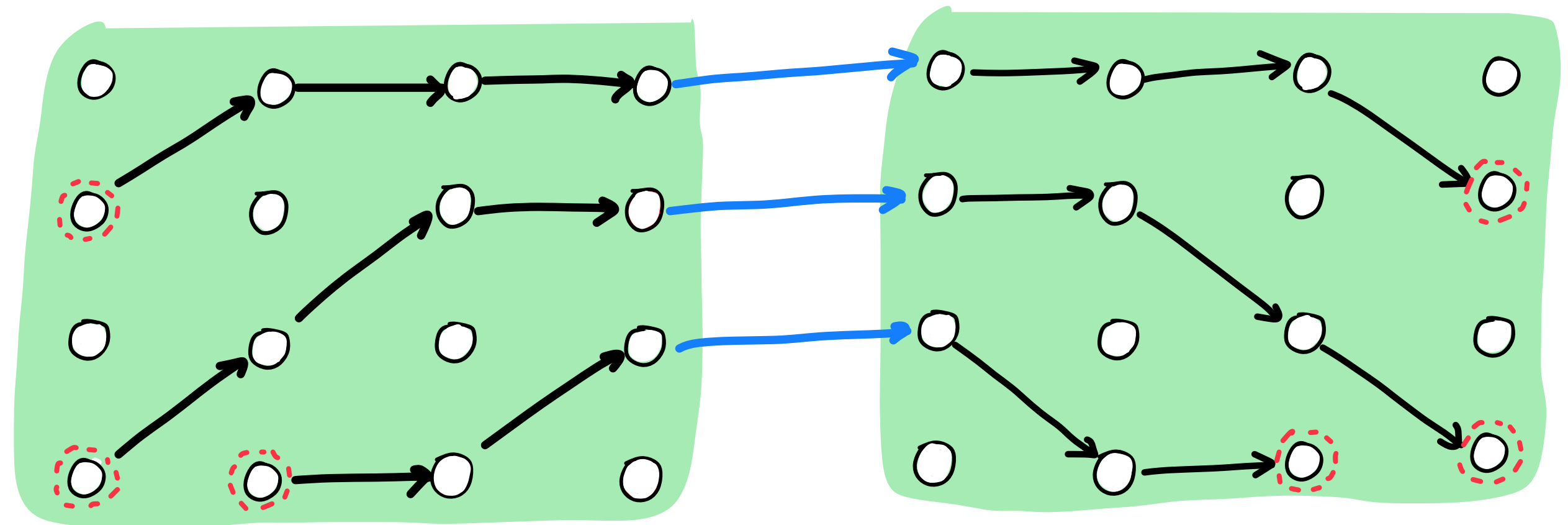
- Only need EITHER ( $SoPL$ ,  $B_{\Sigma^1_1}$  Pigeon)  $\leq$  END OF POTENTIAL LINE
- can we do  $SoPL \leq$  END OF POTENTIAL LINE?

where to put the source?



$SoPL$

or ?



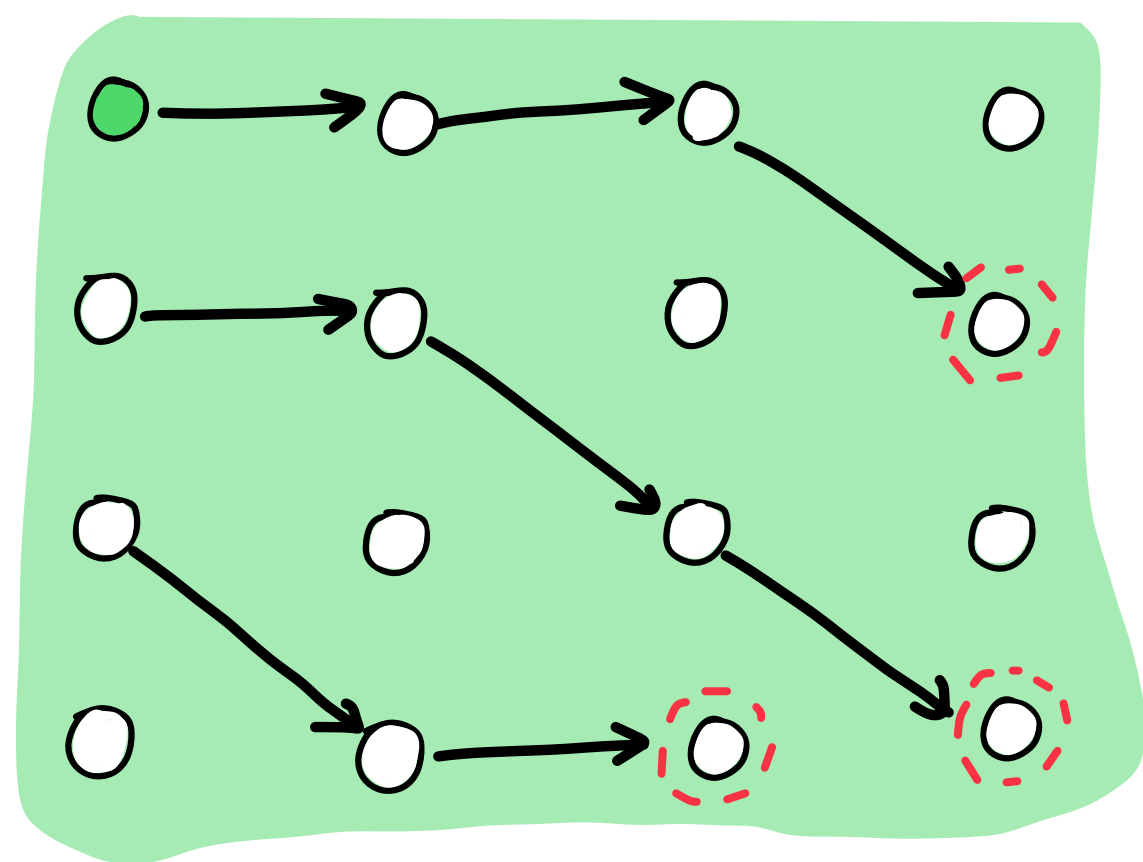
reversed  $SoPL$

$SoPL$

END OF POTENTIAL LINE

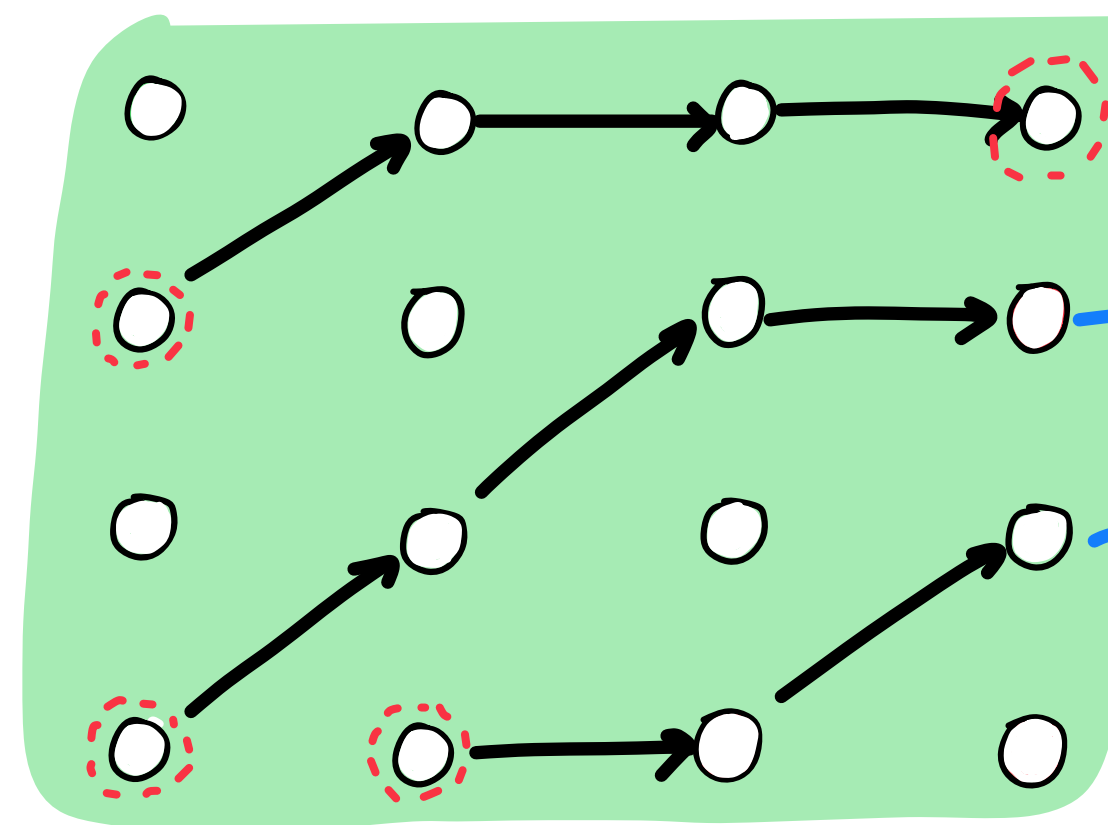
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- Only need EITHER ( $SoPL$ ,  $B_{\exists} \text{ Pigeon}$ )  $\leq$  END OF POTENTIAL LINE
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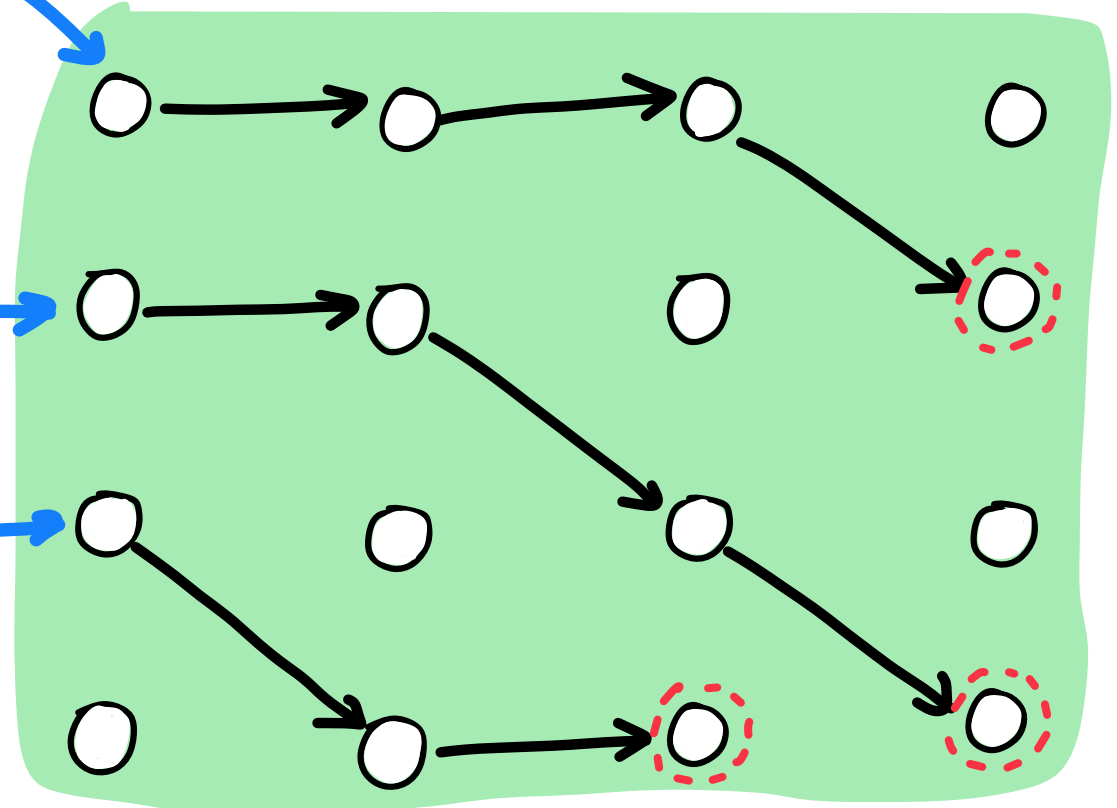
$SoPL$

or ?



reversed  $SoPL$

introduces a bad solution...



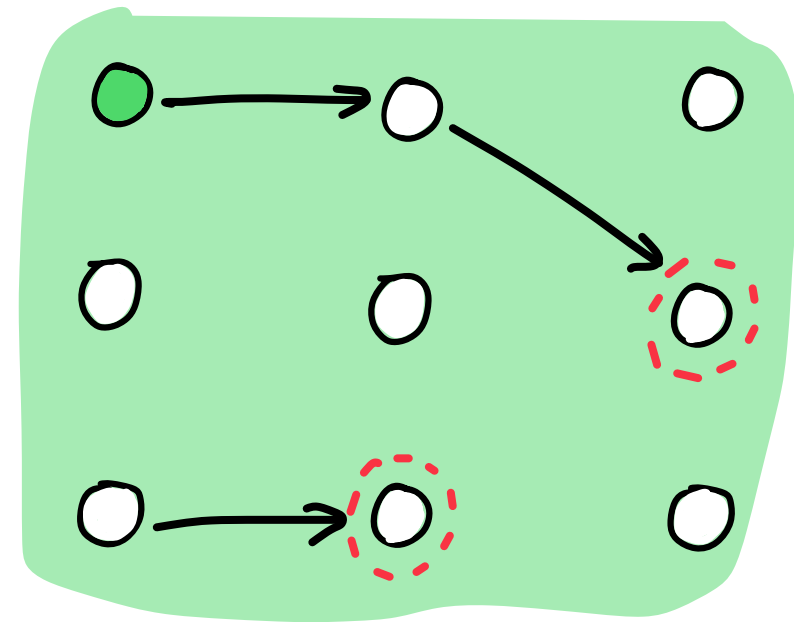
$SoPL$

END OF POTENTIAL LINE

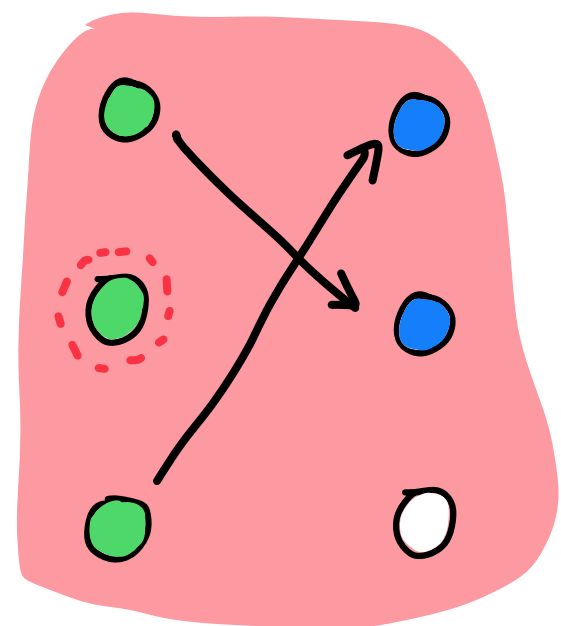
# More on $E_{OPL} = PLS \cap PPAD$

Idea 3: Use the  $BIM$  Pigeon instance to incorporate the source

SoPL

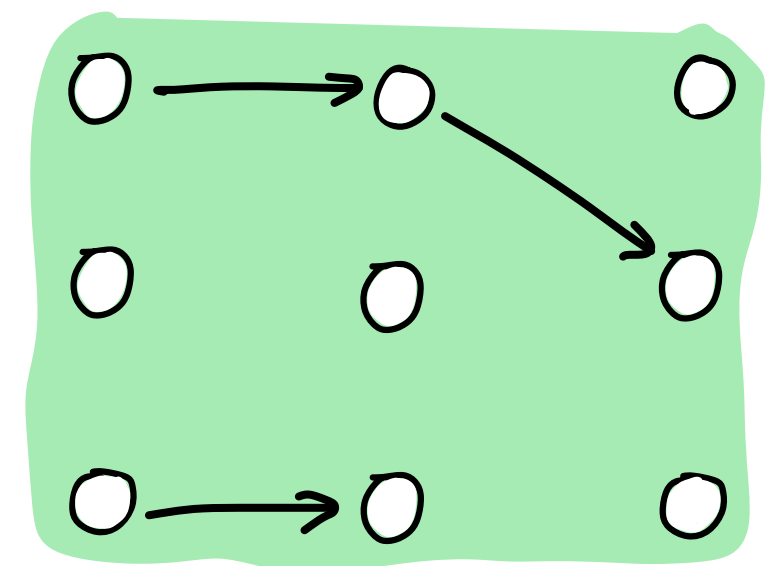
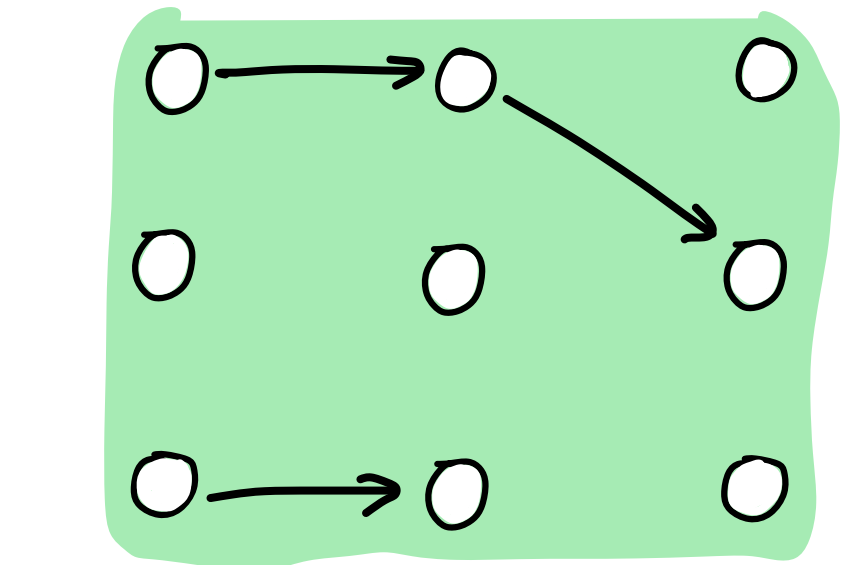
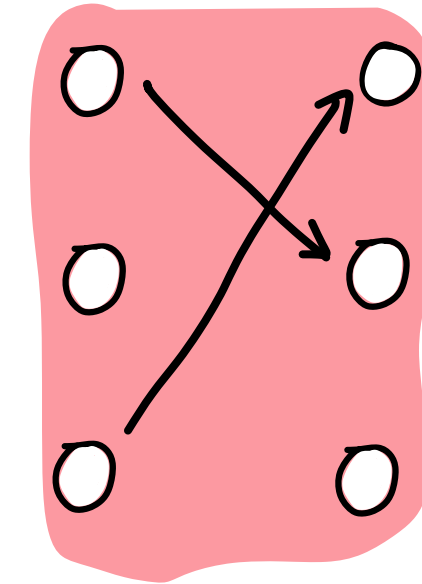
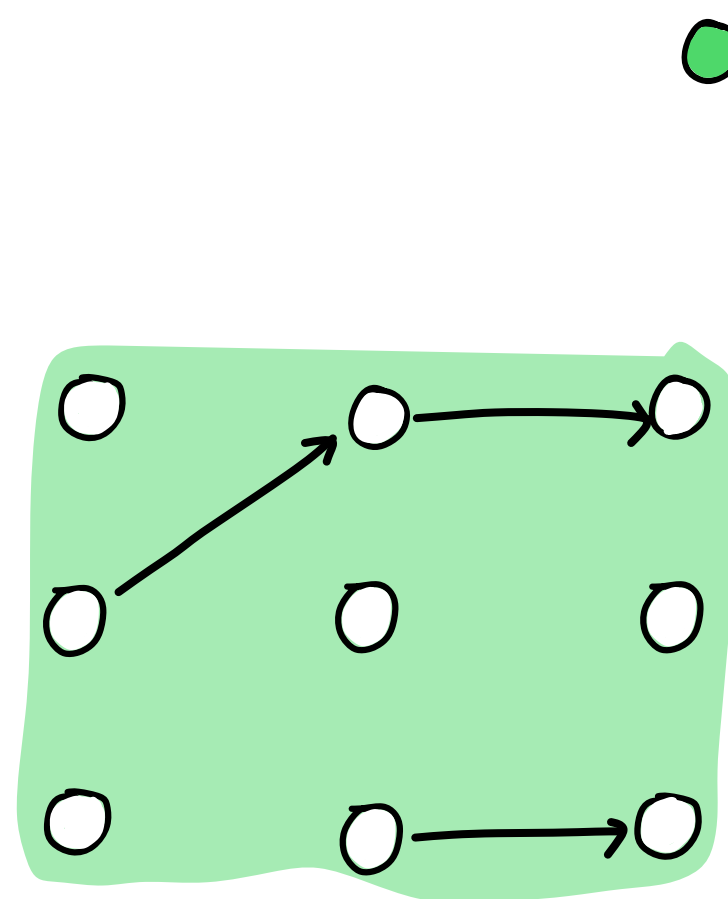


$\wedge$



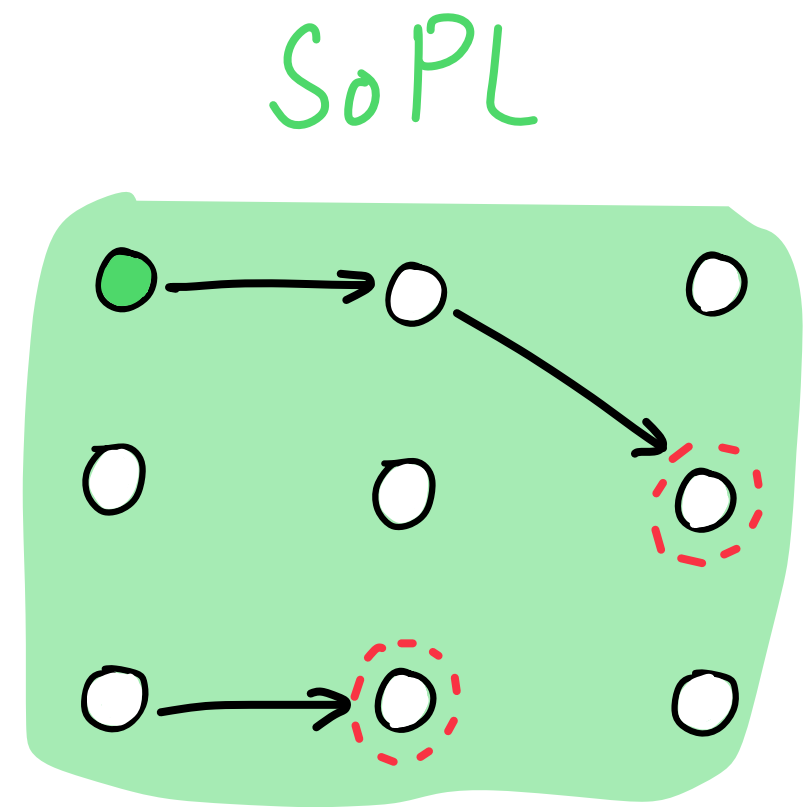
$BIM$  Pigeon

$\Rightarrow$

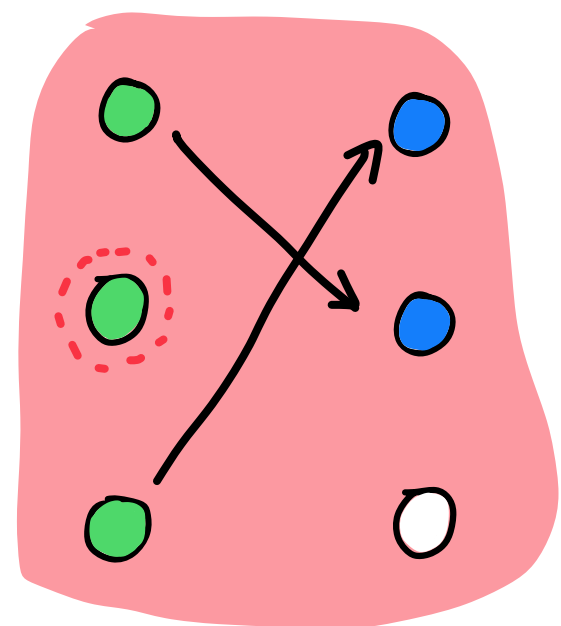


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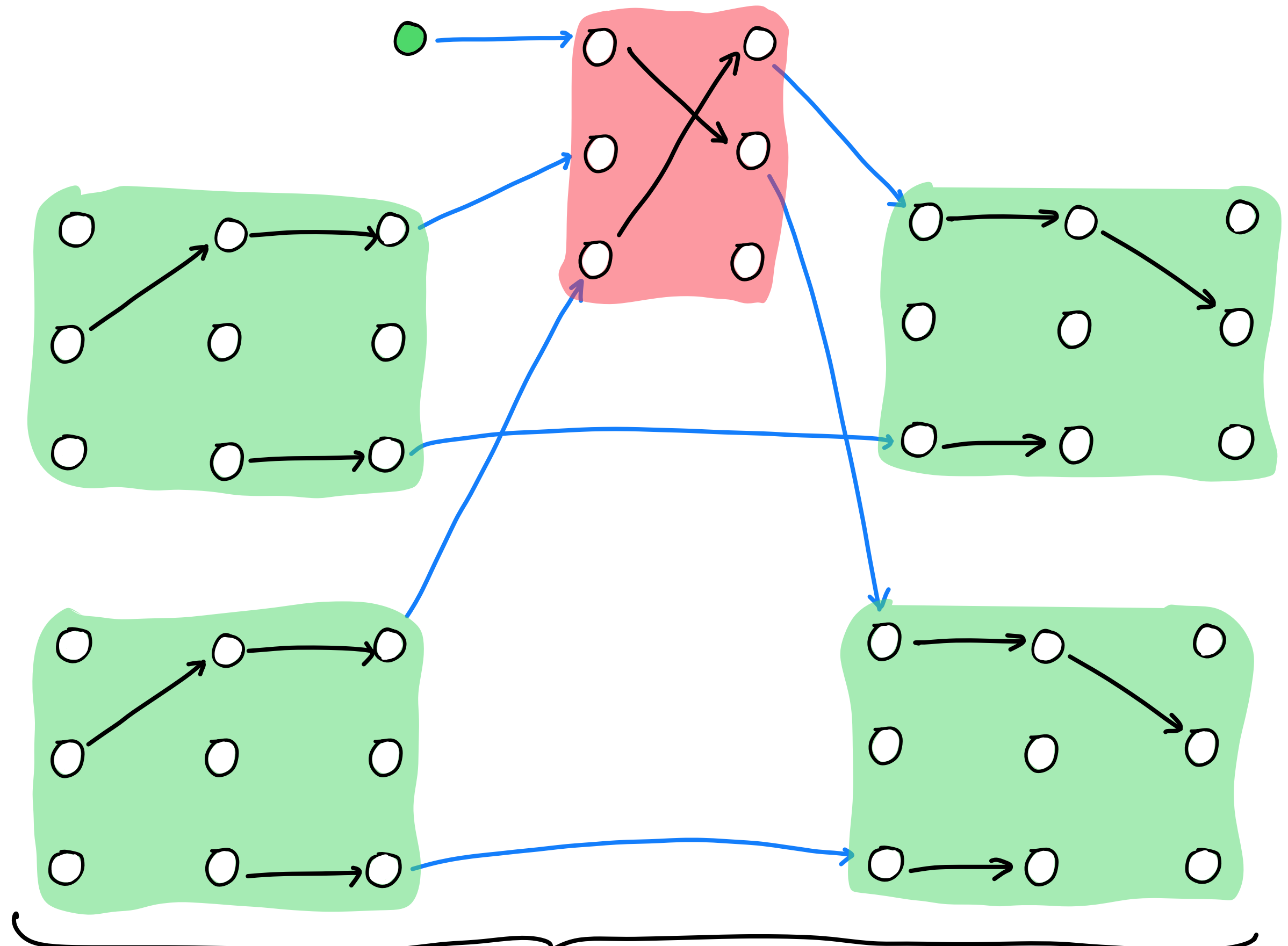


$\wedge$



$BIM$  Pigeon

$\cong$

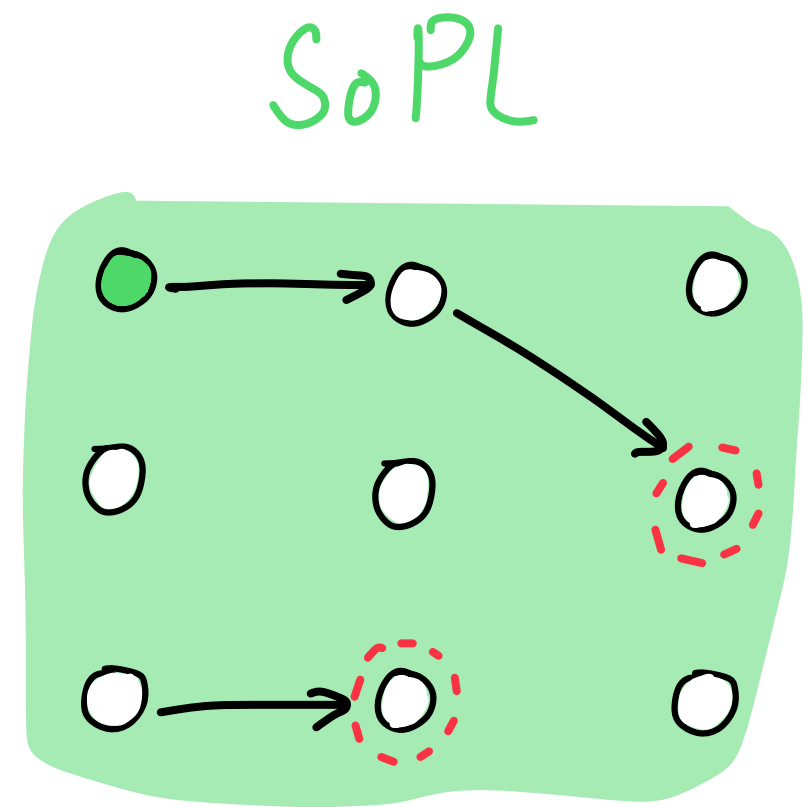


$EoPL$  instance

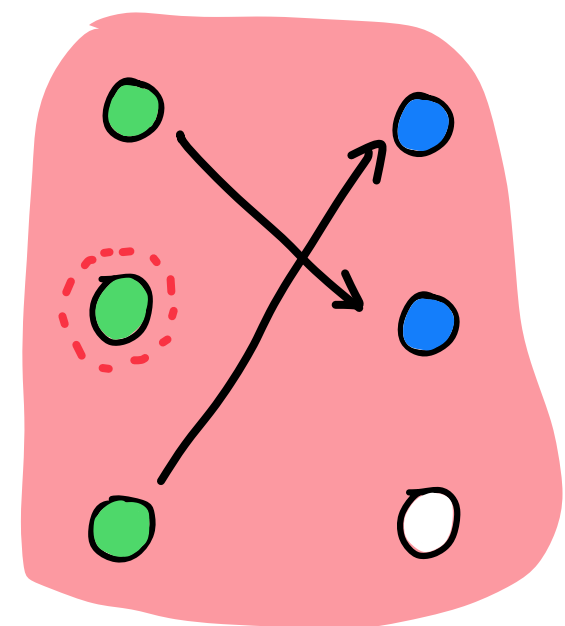


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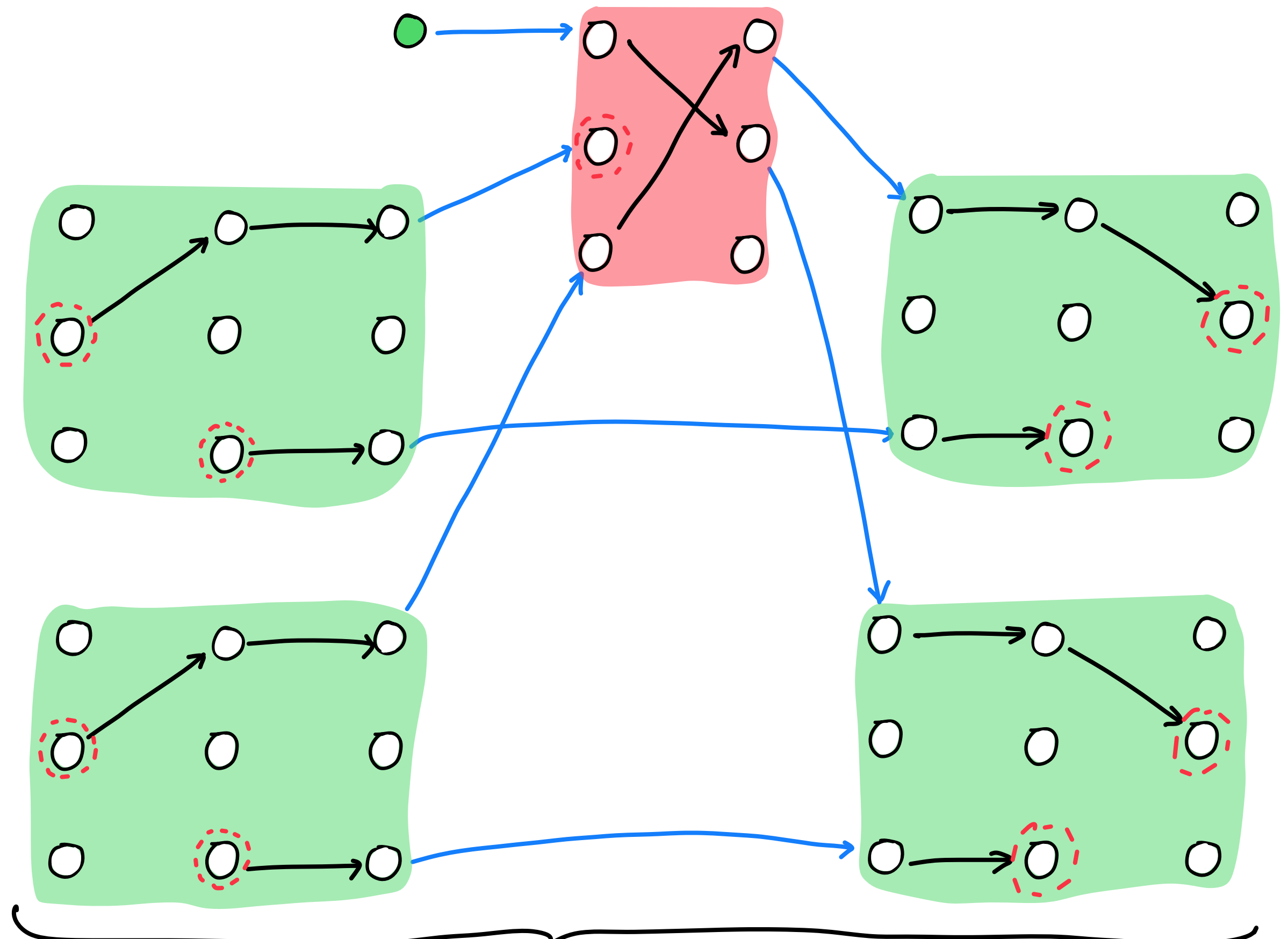


$\wedge$



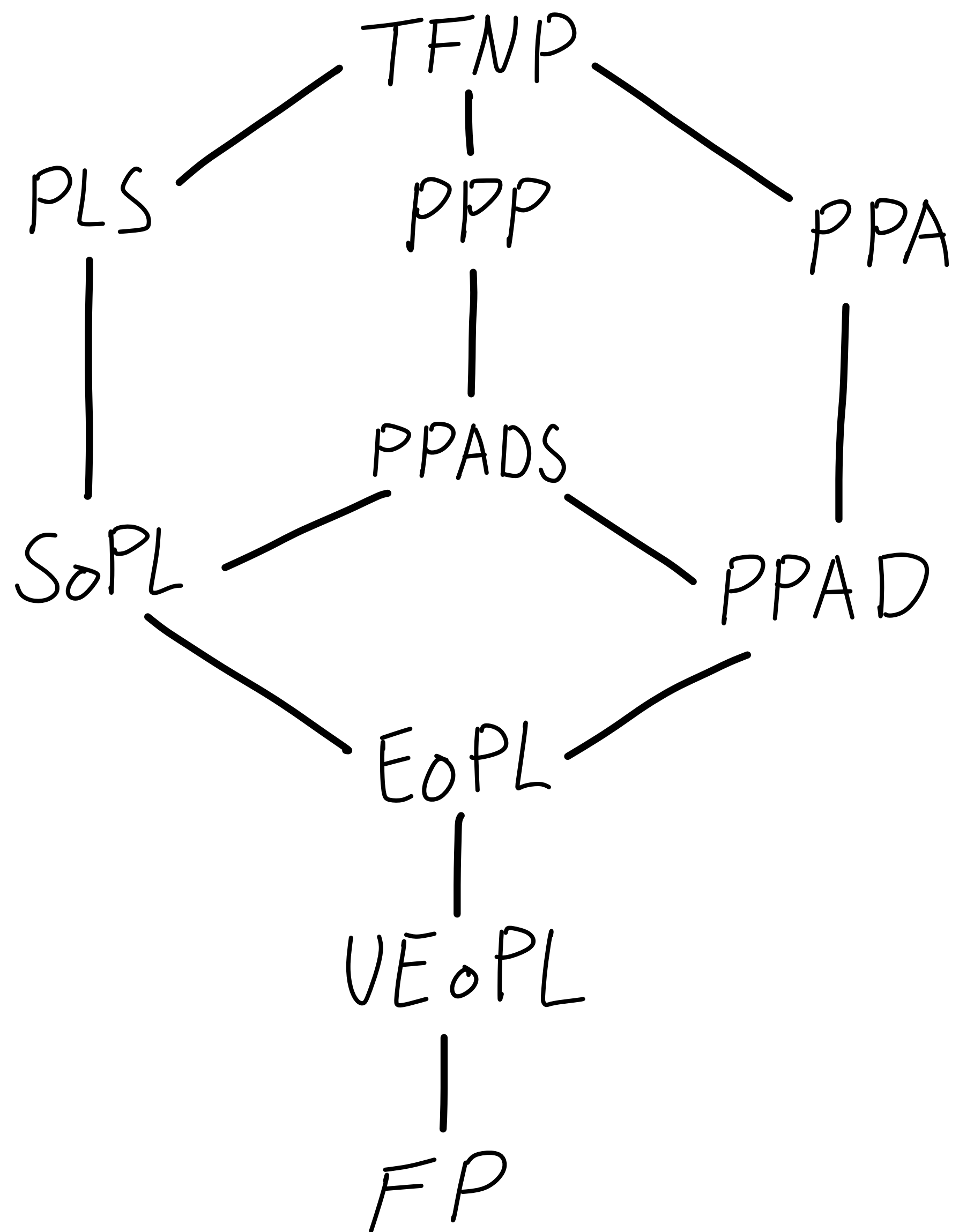
$BIM$  Pigeon

$\cong$



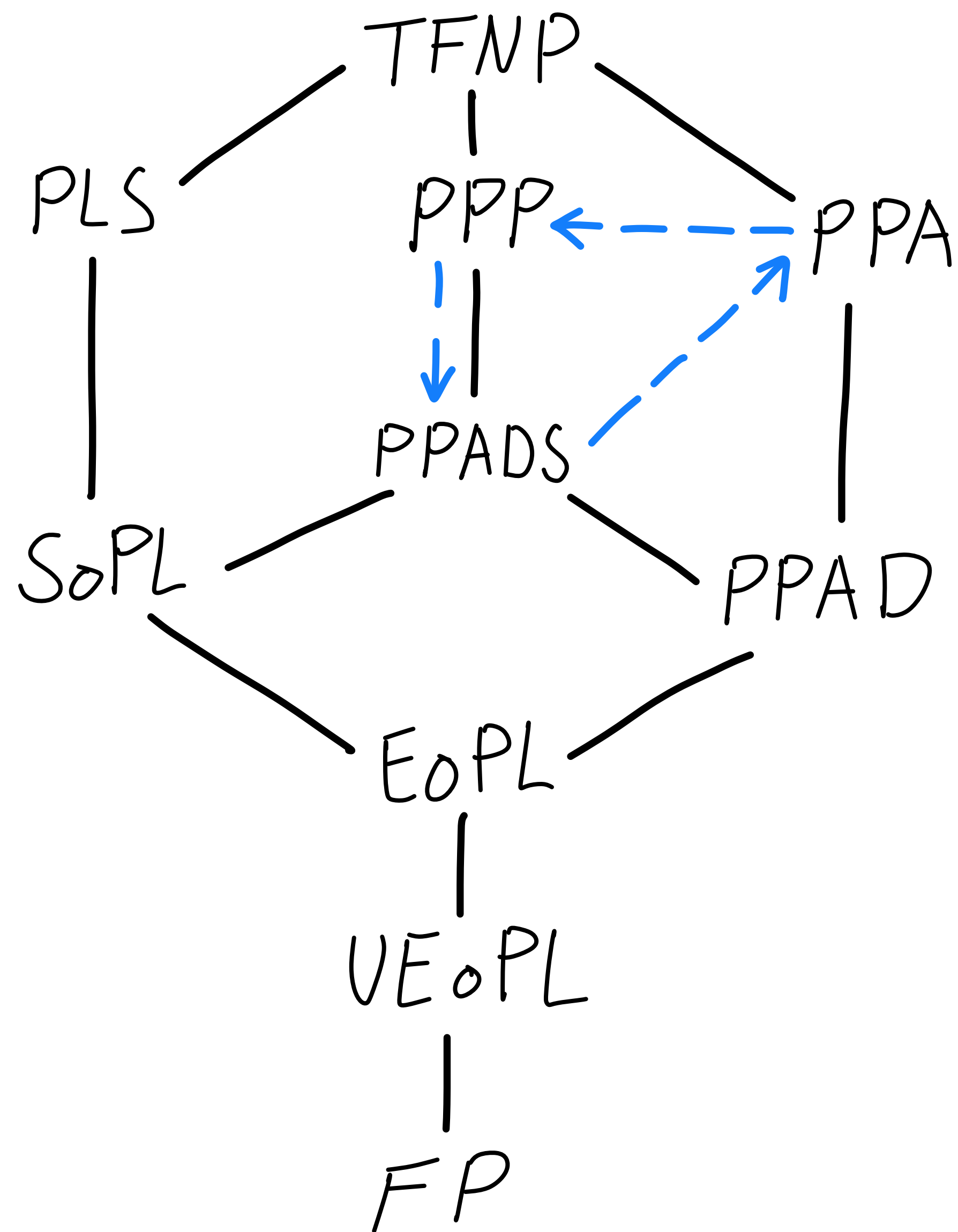
$EoPL$  instance

Result II: no further black-box collapses



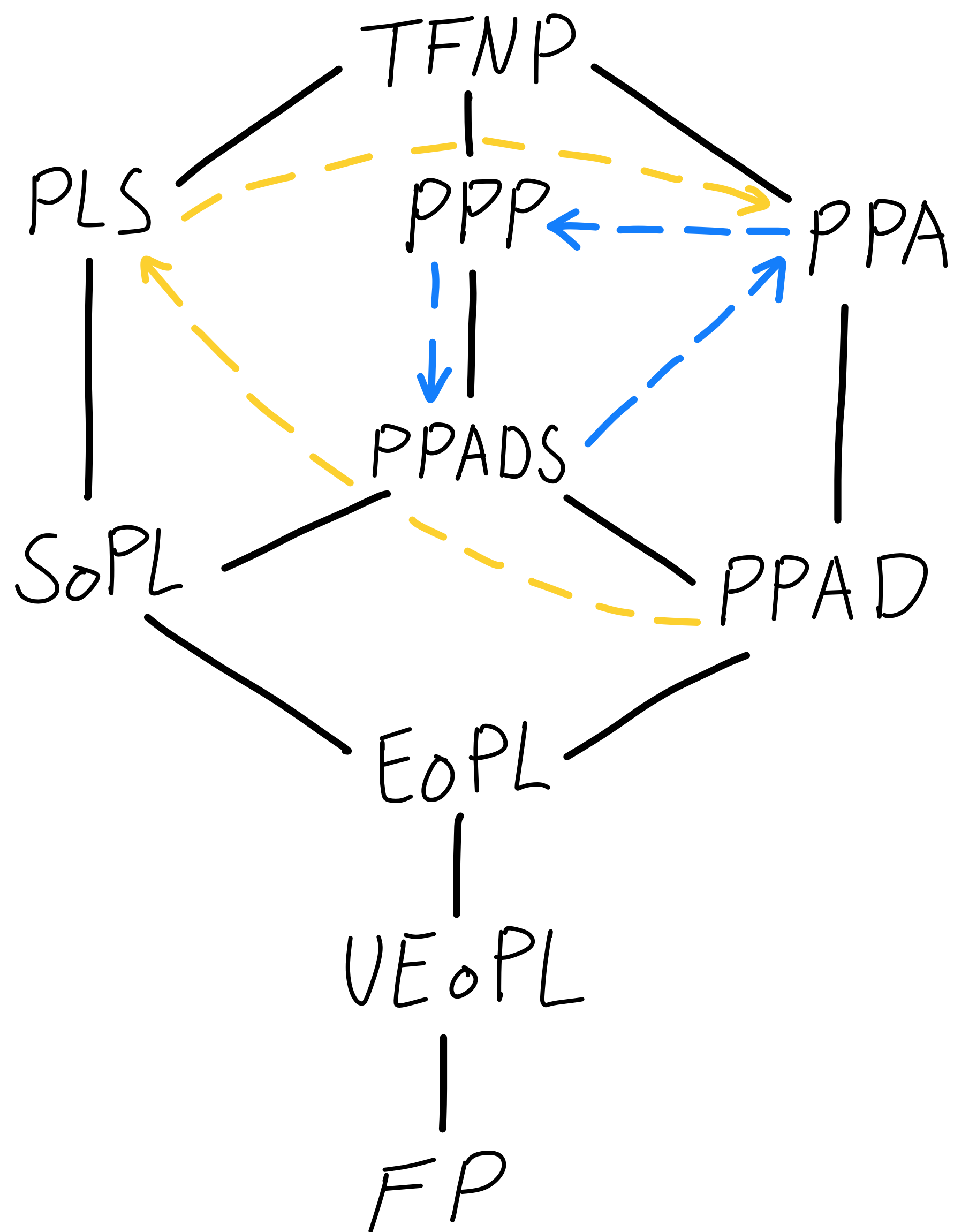
- unconditional separation  $\Rightarrow P \neq NP$
- $CLS$ ,  $SoPL$  and  $EoPL$  collapses were **black-box!**
- **black-box** separation  $\Rightarrow$  oracle separation
- initiated by **Beame et. al. 98'**

# Result II: no further black-box collapses



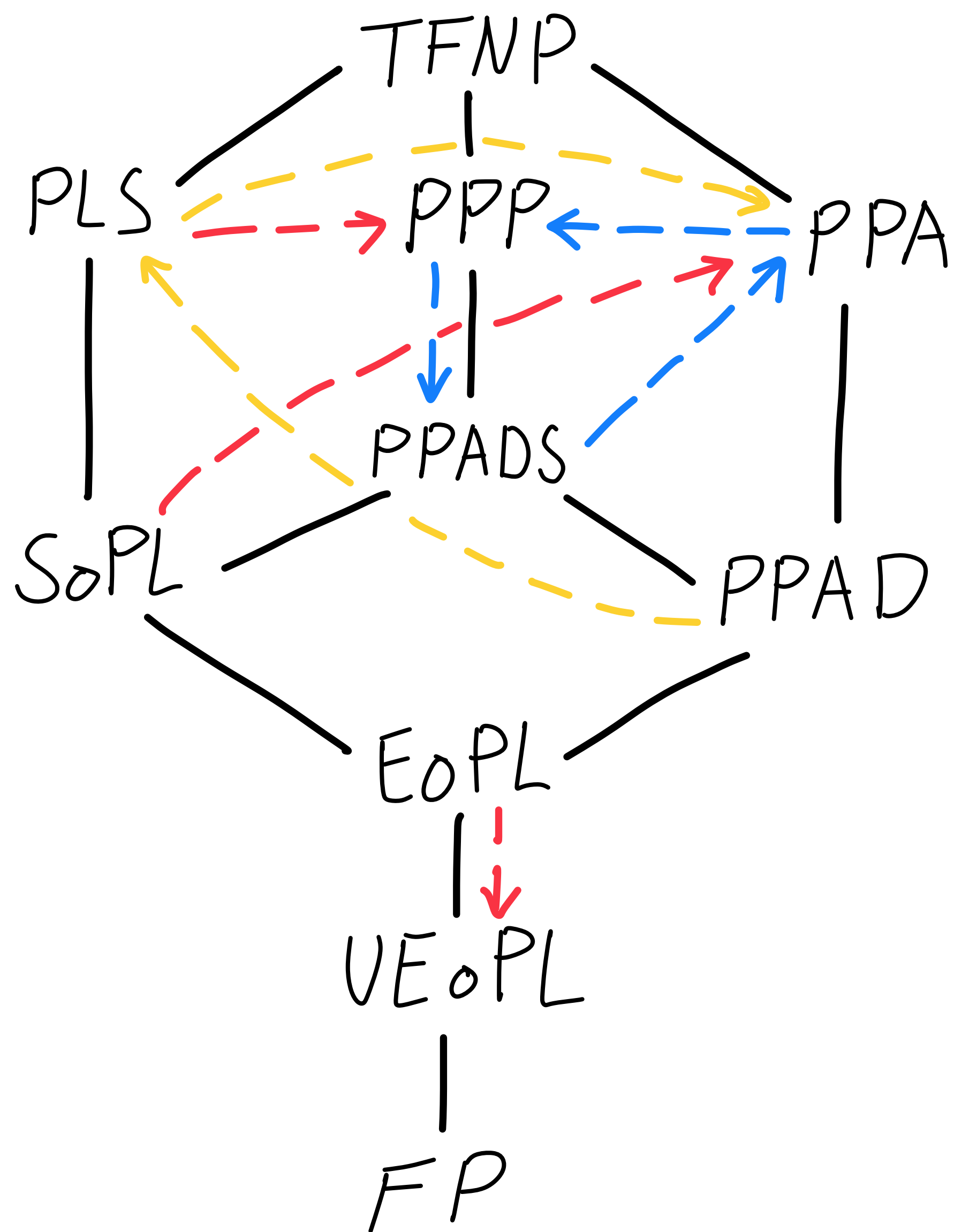
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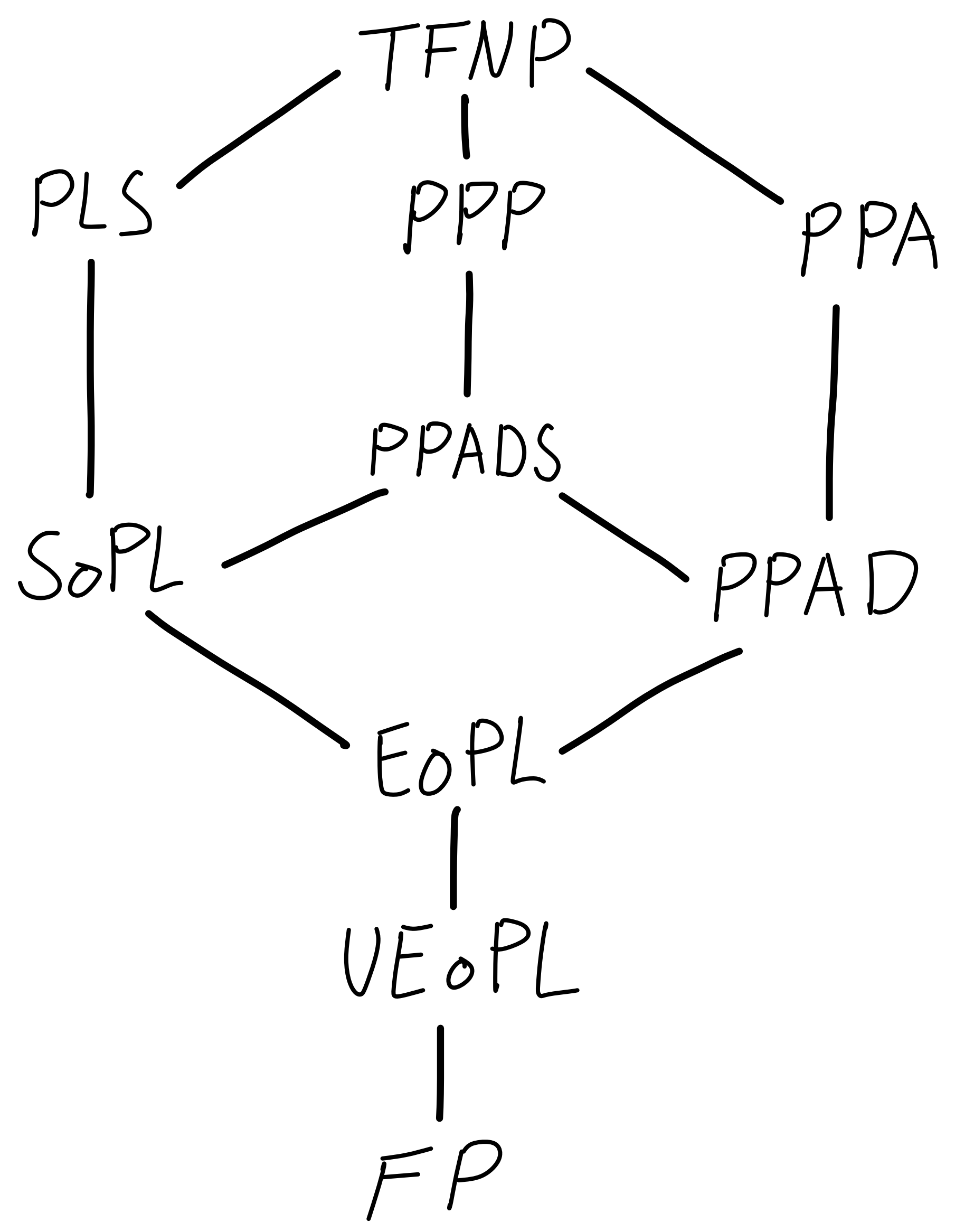
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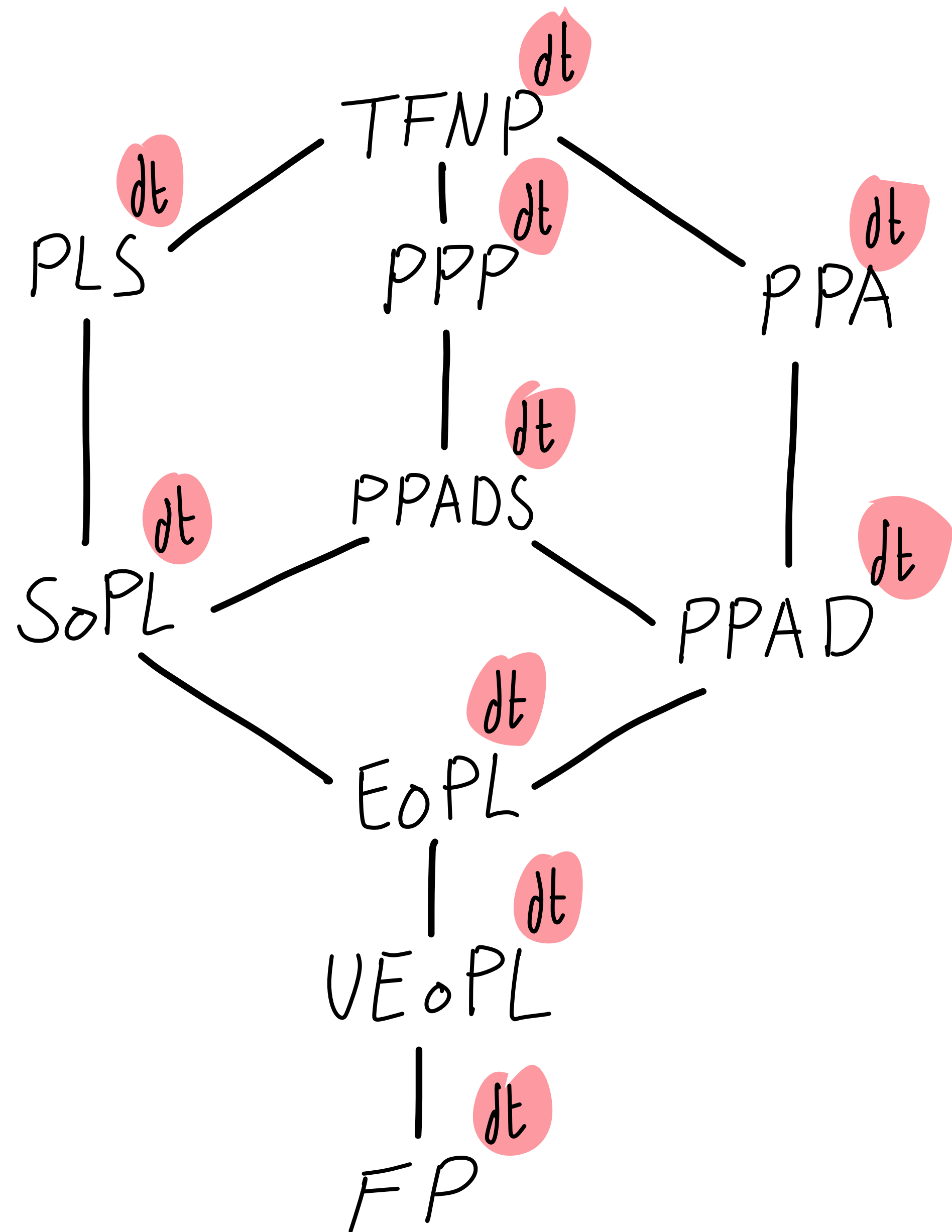


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- 2022

# Tool 1: black-box classes



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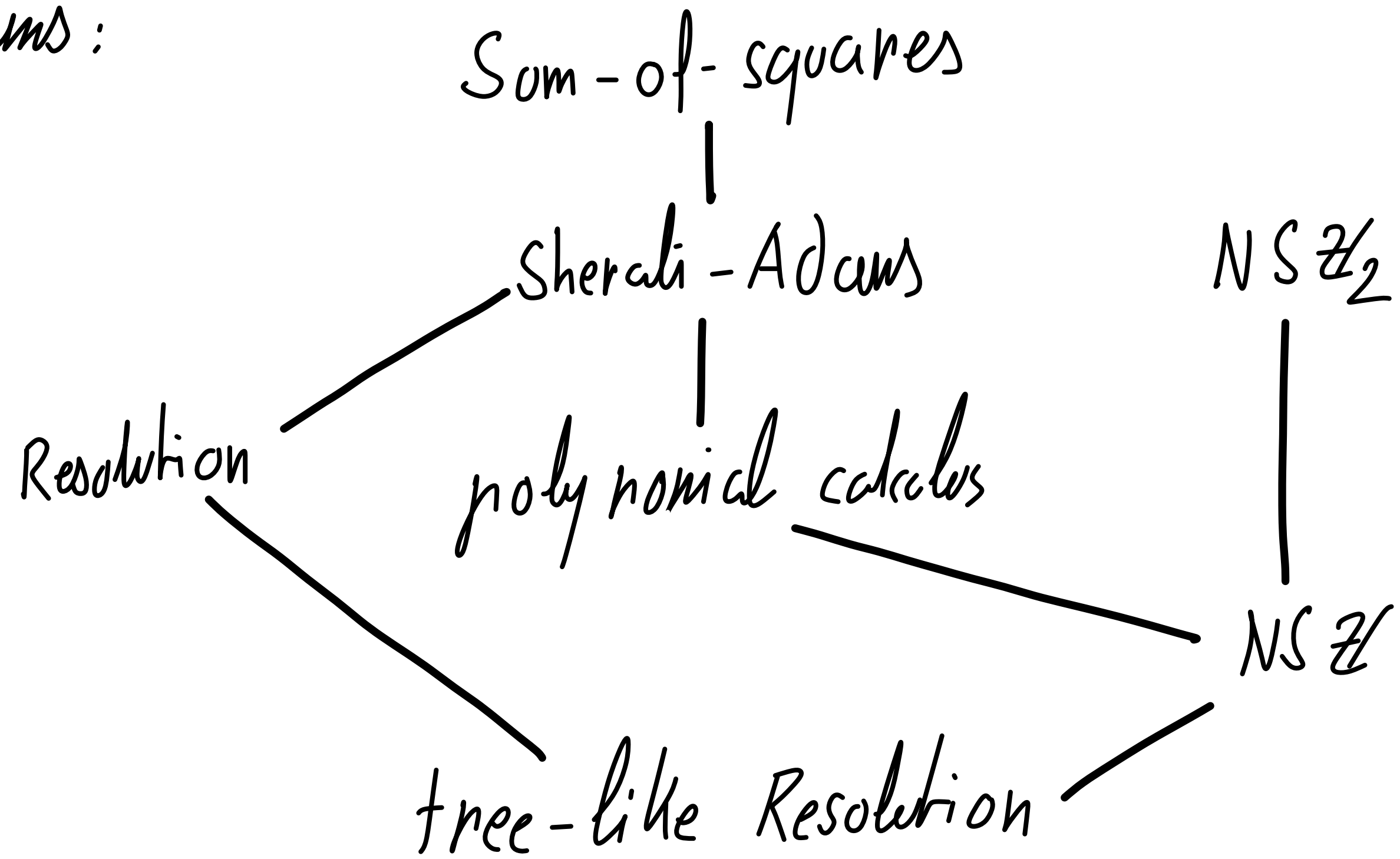


- study "query analogues" a.k.a type-2 complexity
- reduction carried out by shallow decision trees

# Tool 2: proof systems

Is there a short way to show that a CNF is unsat?

- hierarchy of proof systems:





## Tool 2: proof systems

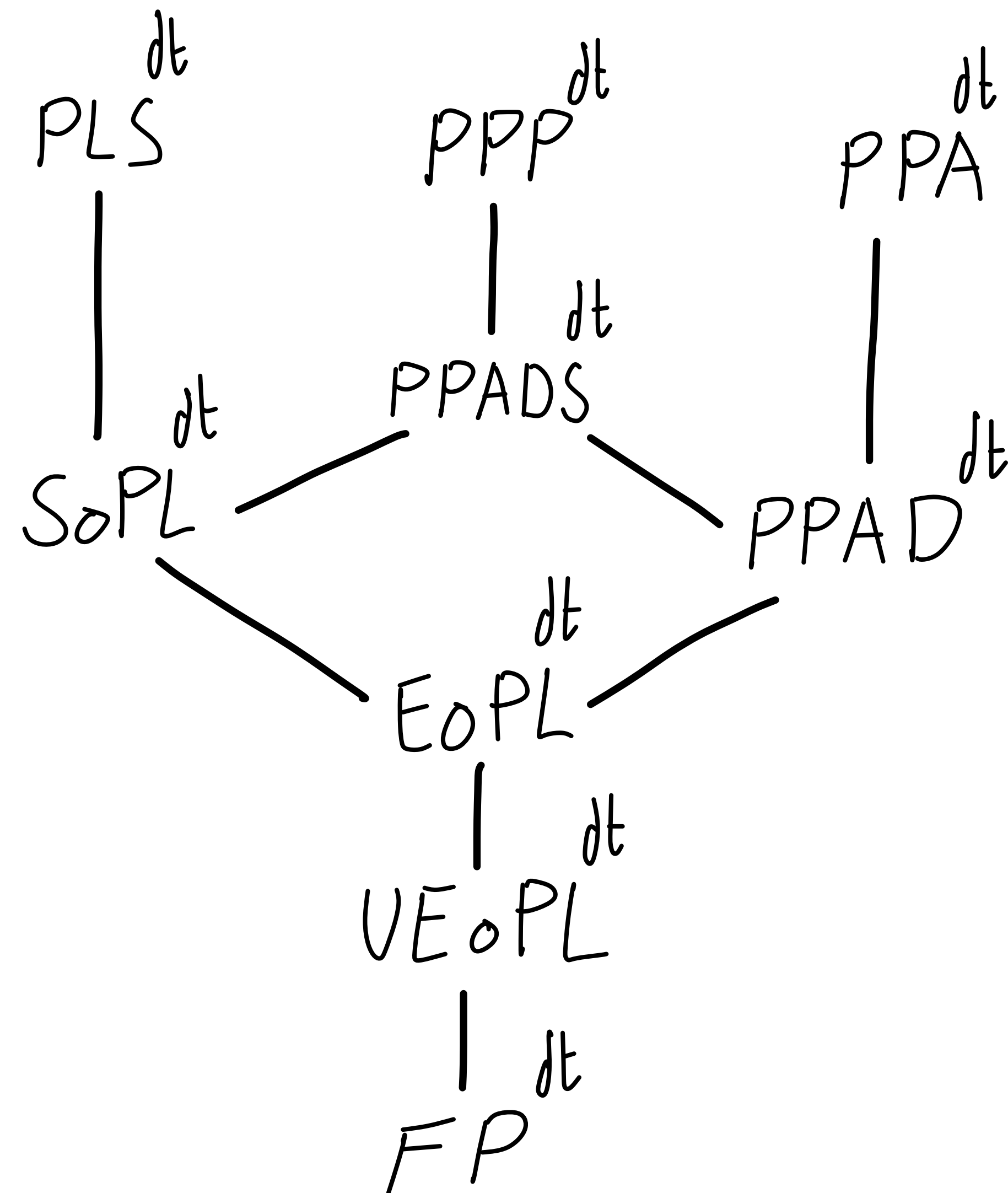
Is there a short way to show that a CNF is unsat?

- hierarchy of proof systems
- $TFNP^{dt}$  search problems can be translated to CNF contradictions

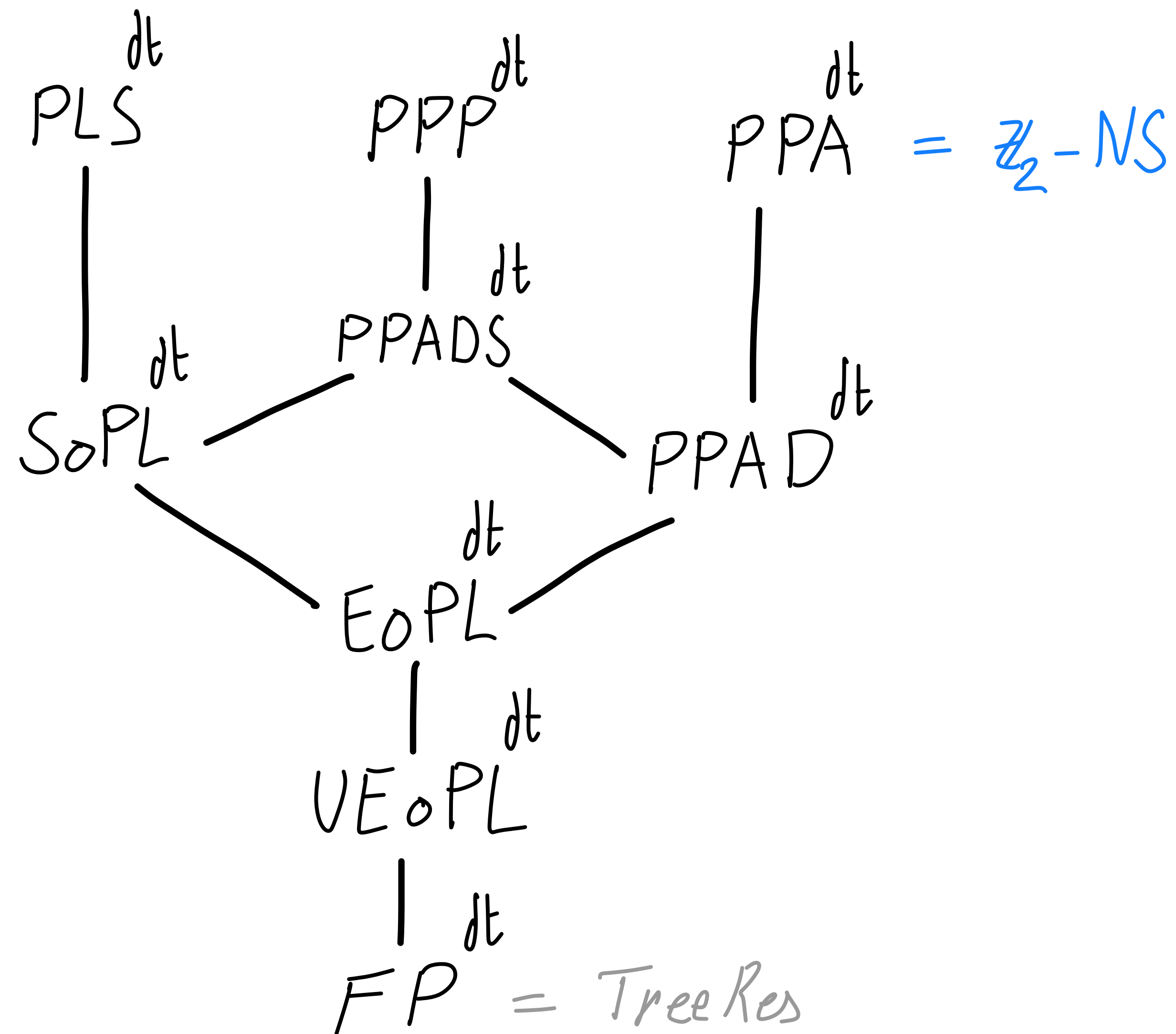
LONELY  $\Rightarrow$

"there is **no** unmatched vertex in this odd-size graph"

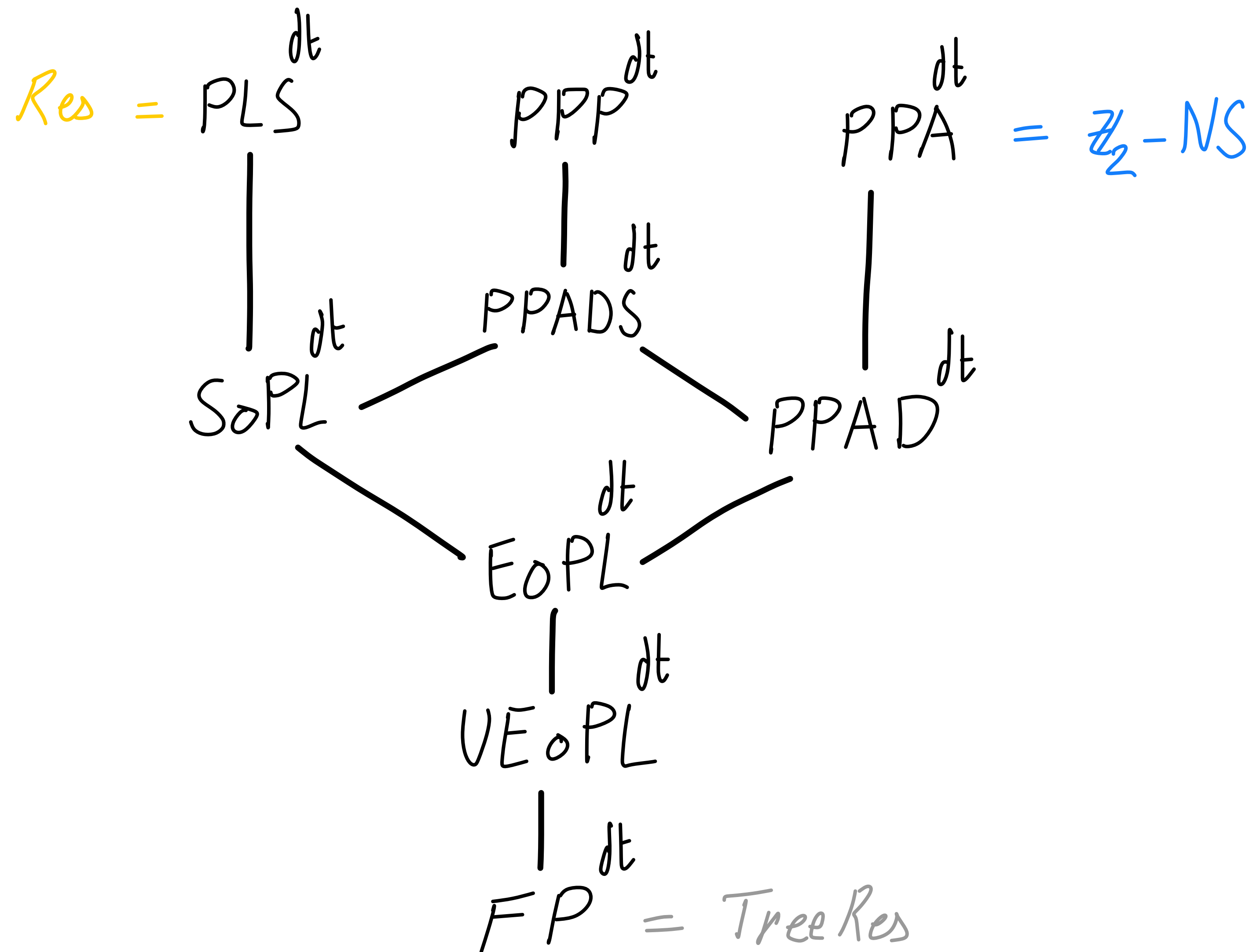
# Tool 3: characterizations



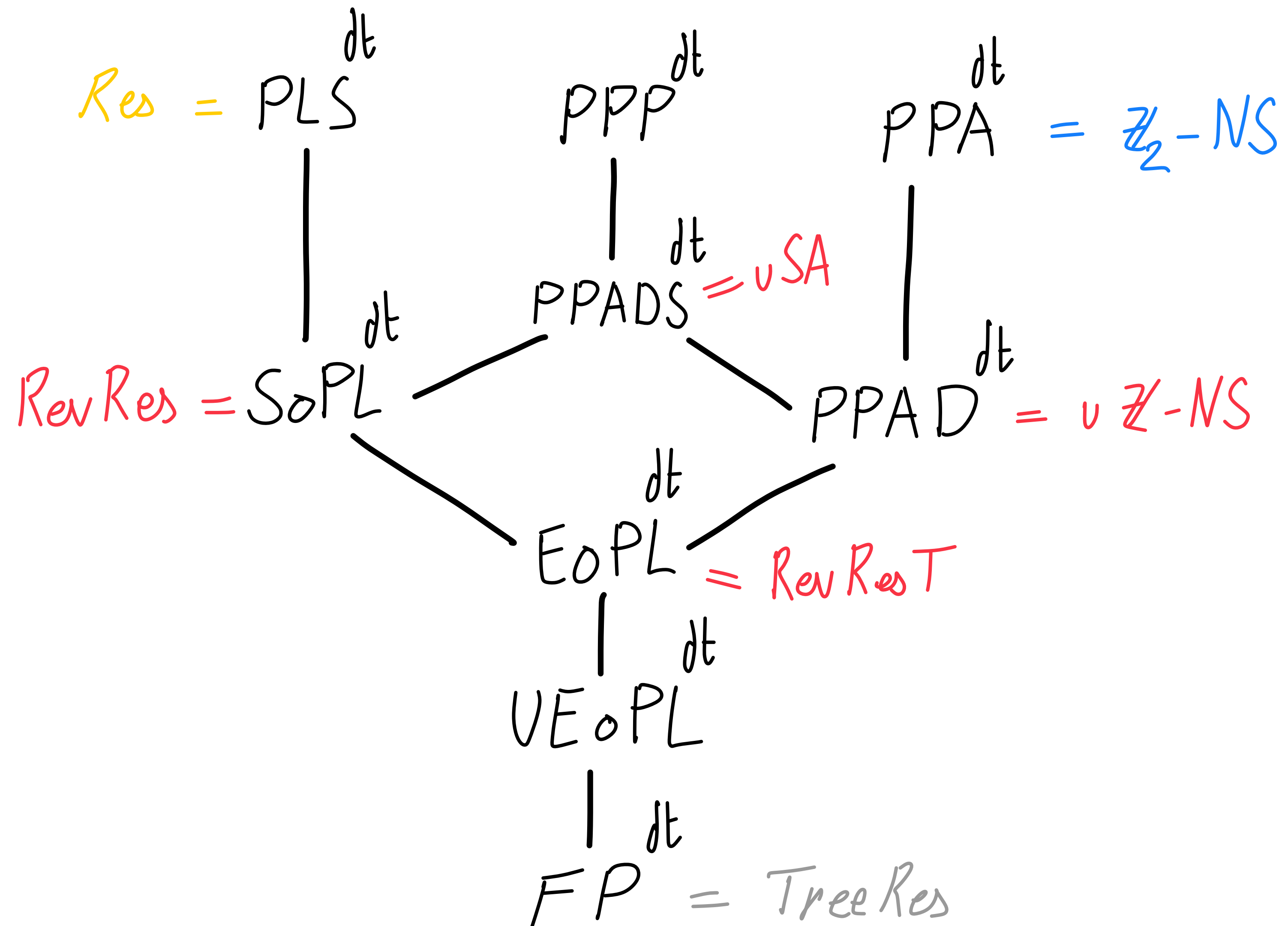
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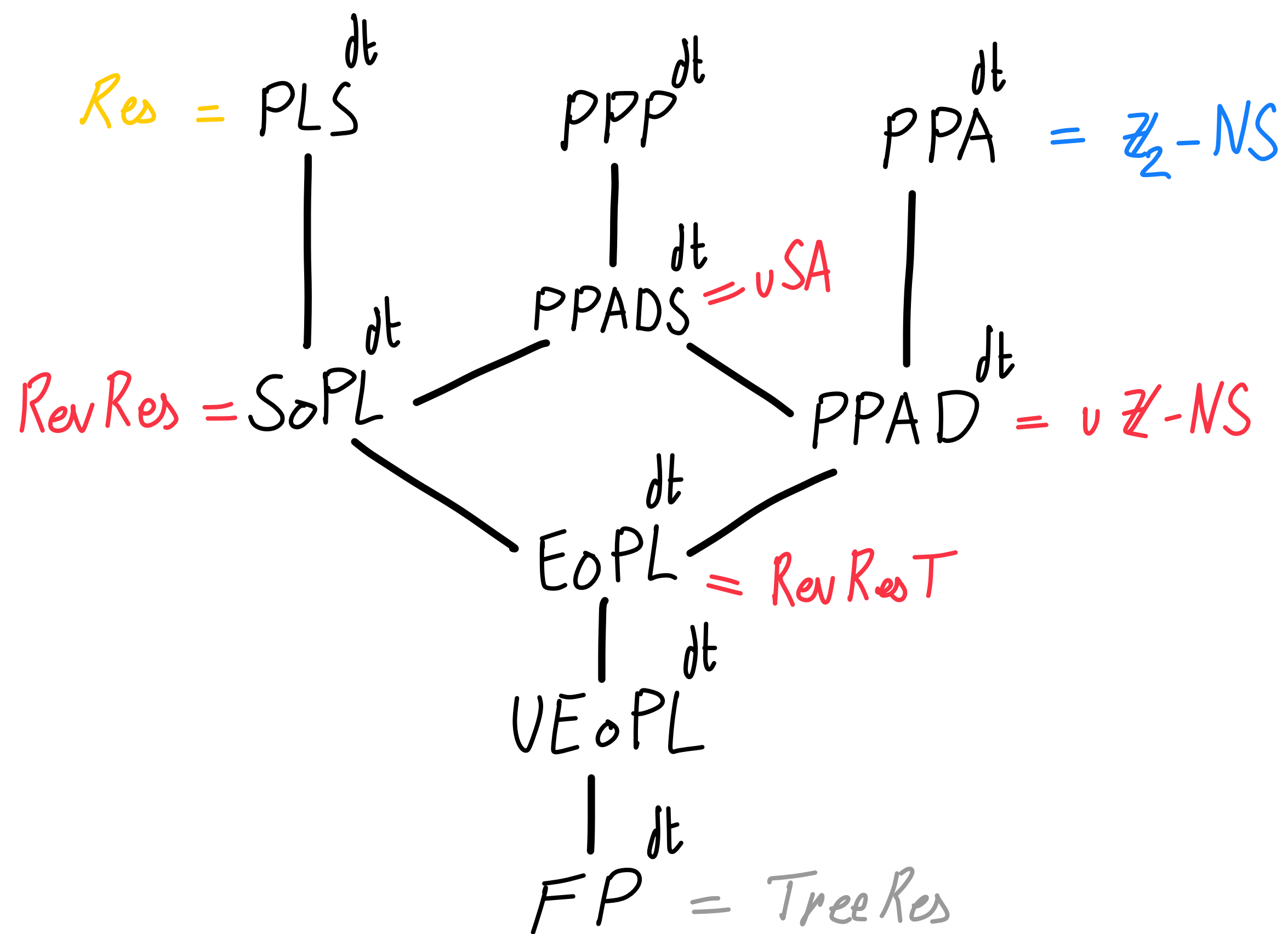
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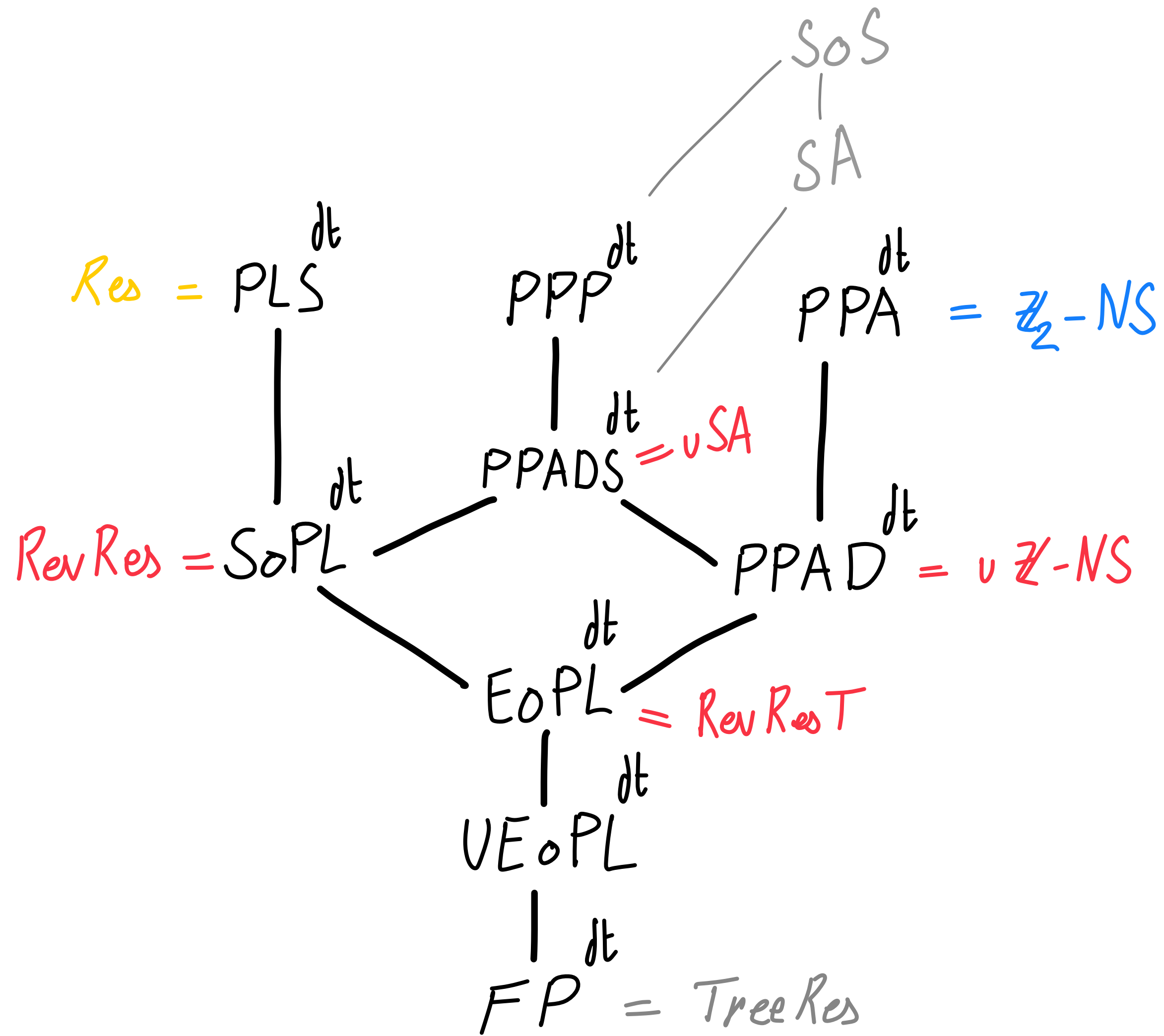


# Tool 3: characterizations



- NEW VIEW: separate proof systems
- for  $PPP^{dt}$  and  $UEoPL^{dt}$ : further combinatorial insights

# Open problems



## on characterizations

- a system for PPP?
- a system for VEOPL?
- a class for polynomial calculus?

## on the structure of TFNP

- where lies PWPP? Ramsey?  
sonflower?